# Mathematics of Data: From Theory to Computation

Prof. Volkan Cevher volkan.cevher@epfl.ch

Lecture 1: Objects in space

Laboratory for Information and Inference Systems (LIONS) École Polytechnique Fédérale de Lausanne (EPFL)

### EE-556 (Fall 2015)





# License Information for Mathematics of Data Slides

- This work is released under a <u>Creative Commons License</u> with the following terms:
- Attribution
  - The licensor permits others to copy, distribute, display, and perform the work. In return, licensees must give the original authors credit.
- Non-Commercial
  - The licensor permits others to copy, distribute, display, and perform the work. In return, licensees may not use the work for commercial purposes – unless they get the licensor's permission.
- Share Alike
  - The licensor permits others to distribute derivative works only under a license identical to the one that governs the licensor's work.
- Full Text of the License



## Outline

- This class:
  - 1. Linear algebra review
    - Notation
    - Vectors
    - Matrices
    - Tensors
- Next class
  - 1. Review of probability theory



## **Recommended reading material**

- Zico Kolter and Chuong Do, Linear Algebra Review and Reference http://cs229.stanford.edu/section/cs229-linalg.pdf, 2012.
- KC Border, Quick Review of Matrix and Real Linear Algebra http://www.hss.caltech.edu/~kcb/Notes/LinearAlgebra.pdf, 2013.
- Simon Foucart and Holger Rauhut, A mathematical introduction to compressive sensing (Appendix A: Matrix Analysis), Springer, 2013.
- Joel A Tropp, Column subset selection, matrix factorization, and eigenvalue optimization, In Proc. of the 20th Annual ACM-SIAM Symposium on Discrete Algorithms, pp 978–986, SIAM, 2009.



### Motivation

## Motivation

This lecture is intended to help you follow mathematical discussions in data sciences, which rely heavily on basic linear algebra concepts.





## Notation

- ▶ Scalars are denoted by lowercase letters (e.g. k)
- Vectors by lowercase boldface letters (e.g., x)
- Matrices by uppercase boldface letters (e.g. A)
- Component of a vector  $\mathbf{x}$ , matrix  $\mathbf{A}$  as  $x_i$ ,  $a_{ij}$  &  $A_{i,j,k,\ldots}$  respectively.
- Sets by uppercase calligraphic letters (e.g. S).



### Vector spaces

#### Note:

We focus on the field of real numbers  $(\mathbb{R})$  but most of the results can be generalized to the field of complex numbers  $(\mathbb{C})$ .

A vector space or *linear space* (over the field  $\mathbb{R}$ ) consists of

- (a) a set of vectors  $\mathcal{V}$
- (b) an addition operation:  $\mathcal{V} \times \mathcal{V} \rightarrow \mathcal{V}$
- (c) a scalar multiplication operation:  $\mathbb{R} \times \mathcal{V} \rightarrow \mathcal{V}$
- (d) a distinguished element  $\mathbf{0}\in\mathcal{V}$

and satisfies the following properties:

1. 
$$\mathbf{x} + \mathbf{y} = \mathbf{y} + \mathbf{x}$$
,  $\forall \mathbf{x}, \mathbf{y} \in \mathcal{V}$   
2.  $(\mathbf{x} + \mathbf{y}) + \mathbf{z} = \mathbf{x} + (\mathbf{y} + \mathbf{z}), \forall \mathbf{x}, \mathbf{y}, \mathbf{z} \in \mathcal{V}$   
3.  $\mathbf{0} + \mathbf{x} = \mathbf{x}, \forall \mathbf{x} \in \mathcal{V}$   
4.  $\forall \mathbf{x} \in \mathcal{V} \exists (-\mathbf{x}) \in \mathcal{V}$  such that  $\mathbf{x} + (-\mathbf{x}) = \mathbf{0}$   
5.  $(\alpha\beta)\mathbf{x} = \alpha(\beta\mathbf{x}), \quad \forall \alpha, \beta \in \mathbb{R} \quad \forall \mathbf{x} \in \mathcal{V}$   
6.  $\alpha(\mathbf{x} + \mathbf{y}) = \alpha\mathbf{x} + \alpha\mathbf{y}, \quad \forall \alpha \in \mathbb{R} \quad \forall \mathbf{x}, \mathbf{y} \in \mathcal{V}$   
7.  $\mathbf{1x} = \mathbf{x}, \forall \mathbf{x} \in \mathcal{V}$ 

commutative under addition associative under addition 0 being additive identity -x being additive inverse associative under scalar multiplication distributive 1 being multiplicative identity





### Vector spaces contd.

### Example (Vector space)

1.  $\mathcal{V}_1 = \{\mathbf{0}\}$  for  $\mathbf{0} \in \mathbb{R}^p$ 2.  $\mathcal{V}_2 = \mathbb{R}^p$ 3.  $\mathcal{V}_3 = \sum_{i=1}^k \alpha_i \mathbf{x}_i$  for  $\alpha_i \in \mathbb{R}$  and  $\mathbf{x}_i \in \mathbb{R}^p$ 

It is straight forward to show that  $V_1$ ,  $V_2$ , and  $V_3$  satisfy properties 1–7 above.

#### Definition (Subspace)

A subspace is a vector space that is a *subset* of another vector space.

### Example (Subspace)

 $\mathcal{V}_1$ ,  $\mathcal{V}_2$ , and  $\mathcal{V}_3$  in the example above are subspaces of  $\mathbb{R}^p$ 



## Vector spaces contd.

# Definition (Span)

The span of a set of vectors,  $\{x_1,x_2,\ldots,x_k\}$ , is the set of all possible linear combinations of these vectors; i.e.,

span { $\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_k$ } = { $\alpha_1 \mathbf{x}_1 + \alpha_2 \mathbf{x}_2 + \dots + \alpha_k \mathbf{x}_k \mid \alpha_1, \alpha_2, \dots, \alpha_k \in \mathbb{R}$ }.

### Definition (Linear independence)

A set of vectors,  $\{x_1, x_2, \dots, x_k\}$ , is linearly independent if

 $\alpha_1 \mathbf{x}_1 + \alpha_2 \mathbf{x}_2 + \dots + \alpha_k \mathbf{x}_k = \mathbf{0} \implies \alpha_1 = \alpha_2 = \dots = \alpha_k = \mathbf{0}.$ 

## Definition (Basis)

The basis of a vector space,  $\mathcal V,$  is a set of vectors  $\{\mathbf x_1, \mathbf x_2, \ldots, \mathbf x_k\}$  that satisfy (a)  $\mathcal V = \mathrm{span}\,\{\mathbf x_1, \mathbf x_2, \ldots, \mathbf x_k\},$  (b)  $\{\mathbf x_1, \mathbf x_2, \ldots, \mathbf x_k\}$  are linearly independent.

### Definition (Dimension\*)

The dimension of a vector space,  $\mathcal{V}$ , (denoted  $\dim(\mathcal{V})$ ) is the number of vectors in the basis of  $\mathcal{V}$ .

\*We will generalize the concept of affine dimension to the *statistical dimension* of convex objects.



## Vector Norms

### Definition (Vector norm)

A norm of a vector in  $\mathbb{R}^p$  is a function  $\|\cdot\| : \mathbb{R}^p \to \mathbb{R}$  such that for all vectors  $\mathbf{x}, \mathbf{y} \in \mathbb{R}^p$  and scalar  $\lambda \in \mathbb{R}$ (a)  $\|\mathbf{x}\| \ge 0$  for all  $\mathbf{x} \in \mathbb{R}^p$  nonnegativity (b)  $\|\mathbf{x}\| = 0$  if and only if  $\mathbf{x} = \mathbf{0}$  definitiveness (c)  $\|\lambda \mathbf{x}\| = |\lambda| \|\mathbf{x}\|$  homogeniety (d)  $\|\mathbf{x} + \mathbf{y}\| \le \|\mathbf{x}\| + \|\mathbf{y}\|$  triangle inequality

- There is a family of  $\ell_q$ -norms parameterized by  $q \in [1, \infty]$ ;
- For  $\mathbf{x} \in \mathbb{R}^p$ , the  $\ell_q$ -norm is defined as  $\|\mathbf{x}\|_q := \left(\sum_{i=1}^p |x_i|^q\right)^{1/q}$ .

### Example

- (1)  $\ell_2$ -norm:  $\|\mathbf{x}\|_2 := \sqrt{\sum_{i=1}^p x_i^2}$  (Euclidean norm)
- (2)  $\ell_1$ -norm:  $\|\mathbf{x}\|_1 := \sum_{i=1}^p |x_i|$  (Manhattan norm)
- (3)  $\ell_{\infty}$ -norm:  $\|\mathbf{x}\|_{\infty} := \max_{i=1,\dots,p} |x_i|$  (Chebyshev norm)





# Definition (Quasi-norm)

A quasi-norm satisfies all the norm properties except (d) triangle inequality, which is replaced by  $\|\mathbf{x} + \mathbf{y}\| \le c (\|\mathbf{x}\| + \|\mathbf{y}\|)$  for a constant  $c \ge 1$ .

# Definition (Semi(pseudo)-norm)

A semi(pseudo)-norm satisfies all the norm properties except (b) definiteness.

### Example

- The  $\ell_q$ -norm is in fact a quasi norm when  $q \in (0, 1)$ , with  $c = 2^{1/q} 1$ .
- ▶ The total variation norm (TV-norm) defined (in 1D):  $\|\mathbf{x}\|_{TV} := \sum_{i=1}^{p-1} |x_{i+1} - x_i|$  is a semi-norm since it fails to satisfy (b); e.g. any  $\mathbf{x} = c(1, 1, ..., 1)^T$  for  $c \neq 0$  will have  $\|\mathbf{x}\|_{TV} = 0$  even though  $\mathbf{x} \neq \mathbf{0}$ .

### Definition ( $\ell_0$ -"norm")

$$\|\mathbf{x}\|_0 = \lim_{q \to 0} \|\mathbf{x}\|_q^q = |\{i : x_i \neq 0\}|$$

The  $\ell_0$ -norm counts the non-zero components of  $\mathbf{x}$ . It is not a norm – it does not satisfy the property (c)  $\Rightarrow$  it is also neither a **quasi**- nor a **semi-norm**.





Problem (s-sparse approximation)					
Find	$\underset{\mathbf{x}\in\mathbb{R}^{p}}{\arg\min} \ \mathbf{x}-\mathbf{y}\ _{2}$	subject to:	$\ \mathbf{x}\ _0 \le s.$		





Problem (*s*-sparse approximation)

Find $\underset{\mathbf{x}\in\mathbb{R}^p}{\arg\min} \ \mathbf{x}-\mathbf{y}\ _2$ subject to: $\ \mathbf{x}\ _0 \leq s$ .				
Solution				
$ \text{Define}  \widehat{\mathbf{y}} \in \mathop{\arg\min}_{\mathbf{x} \in \mathbb{R}^p: \ \mathbf{x}\ _0 \leq s} \ \mathbf{x} - \mathbf{y}\ _2^2  \text{ and let } \widehat{\mathcal{S}} = supp\left(\widehat{\mathbf{y}}\right). $				
We now consider an optimization over sets				
$\widehat{\mathcal{S}} \in \operatorname*{argmin}_{\mathcal{S}: \mathcal{S}  \leq s} \ \mathbf{y}_{\mathcal{S}} - \mathbf{y}\ _2^2.$				
$\in \operatorname*{argmax}_{\mathcal{S}: \mathcal{S}  \leq s} \left\{ \ \mathbf{y}\ _2^2 - \ \mathbf{y}_{\mathcal{S}} - \mathbf{y}\ _2^2 \right\}$				
$ \in \underset{\mathcal{S}: \mathcal{S} \leq s}{\arg\max} \left\{ \ \mathbf{y}_{\mathcal{S}}\ _{2}^{2} \right\} = \underset{\mathcal{S}: \mathcal{S} \leq s}{\arg\max} \sum_{i\in\mathcal{S}} \ y_{i}\ ^{2}  (\equiv \text{ modular approximation problem}). $				

Thus, the best *s*-sparse approximation of a vector is a vector with the *s* largest components of the vector in *magnitude*.



Norm balls

Radius r ball in  $\ell_q$ -norm:

$$\mathcal{B}_q(r) = \{ \mathbf{x} \in \mathbb{R}^p : \|\mathbf{x}\|_q \le r \}$$

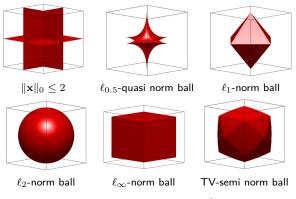


Table: Example norm balls in  $\mathbb{R}^3$ 





### Inner products

#### Definition (Inner product)

The inner product of any two vectors  $\mathbf{x}, \mathbf{y} \in \mathbb{R}^p$  (denoted by  $\langle \cdot, \cdot \rangle$ ) is defined as  $\langle \mathbf{x}, \mathbf{y} \rangle = \mathbf{x}^T \mathbf{y} = \sum_i^p x_i y_i$ .

The inner product satisfies the following properties:

 1.  $\langle \mathbf{x}, \mathbf{y} \rangle = \langle \mathbf{y}, \mathbf{x} \rangle, \forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^p$  symmetry

 2.  $\langle (\alpha \mathbf{x} + \beta \mathbf{y}), \mathbf{z} \rangle = \langle \alpha \mathbf{x}, \mathbf{z} \rangle + \langle \beta \mathbf{y}, \mathbf{z} \rangle, \forall \alpha, \beta \in \mathbb{R}, \forall \mathbf{x}, \mathbf{y}, \mathbf{z} \in \mathbb{R}^p$  linearity

 3.  $\langle \mathbf{x}, \mathbf{x} \rangle \ge 0, \forall \mathbf{x} \in \mathbb{R}^p$  positive definiteness

Important relations involving the inner product:

- Hölder's inequality:  $|\langle \mathbf{x}, \mathbf{y} \rangle| \leq ||\mathbf{x}||_q ||\mathbf{y}||_r$ , where r > 1 and  $\frac{1}{q} + \frac{1}{r} = 1$
- Cauchy-Schwarz is a special case of Hölder's inequality (q = r = 2)

#### Definition (Inner product space)

An inner product space is a vector space endowed with an inner product.



## Definition (Dual norm)

Let  $\|\cdot\|$  be a norm in  $\mathbb{R}^p$ , then the **dual norm** denoted by  $\|\cdot\|^*$  is defined:

$$\|\mathbf{x}\|^* = \sup_{\|\mathbf{y}\| \leq 1} \mathbf{x}^T \mathbf{y}, \text{ for all } \mathbf{x}, \mathbf{y} \in \mathbb{R}^p$$

- ▶ The dual of the *dual norm* is the original (primal) norm, i.e.,  $\|\mathbf{x}\|^{**} = \|\mathbf{x}\|$ .
- Hölder's inequality  $\Rightarrow \|\cdot\|_q$  is a dual norm of  $\|\cdot\|_r$  when  $\frac{1}{q} + \frac{1}{r} = 1$ .

### Example 1

i)  $\|\cdot\|_2$  is dual of  $\|\cdot\|_2$  (i.e.  $\|\cdot\|_2$  is self-dual):  $\sup\{\mathbf{z}^T\mathbf{x} \mid \|\mathbf{x}\|_2 \le 1\} = \|\mathbf{z}\|_2$ . ii)  $\|\cdot\|_1$  is dual of  $\|\cdot\|_\infty$ , (and vice versa):  $\sup\{\mathbf{z}^T\mathbf{x} \mid \|\mathbf{x}\|_\infty \le 1\} = \|\mathbf{z}\|_1$ .

#### Example 2

What is the dual norm of  $\|\cdot\|_q$  for  $q = 1 + 1/\log(p)$ ?





## Definition (Dual norm)

Let  $\|\cdot\|$  be a norm in  $\mathbb{R}^p$ , then the **dual norm** denoted by  $\|\cdot\|^*$  is defined:

$$\|\mathbf{x}\|^* = \sup_{\|\mathbf{y}\| \leq 1} \mathbf{x}^T \mathbf{y}, \text{ for all } \mathbf{x}, \mathbf{y} \in \mathbb{R}^p$$

- ▶ The dual of the *dual norm* is the original (primal) norm, i.e.,  $\|\mathbf{x}\|^{**} = \|\mathbf{x}\|$ .
- Hölder's inequality  $\Rightarrow \|\cdot\|_q$  is a dual norm of  $\|\cdot\|_r$  when  $\frac{1}{q} + \frac{1}{r} = 1$ .

### Example 1

i)  $\|\cdot\|_2$  is dual of  $\|\cdot\|_2$  (i.e.  $\|\cdot\|_2$  is self-dual):  $\sup\{\mathbf{z}^T\mathbf{x} \mid \|\mathbf{x}\|_2 \le 1\} = \|\mathbf{z}\|_2$ . ii)  $\|\cdot\|_1$  is dual of  $\|\cdot\|_\infty$ , (and vice versa):  $\sup\{\mathbf{z}^T\mathbf{x} \mid \|\mathbf{x}\|_\infty \le 1\} = \|\mathbf{z}\|_1$ .

#### Example 2

What is the dual norm of  $\|\cdot\|_q$  for  $q = 1 + 1/\log(p)$ ?

## Solution By Hölder's inequality, $\|\cdot\|_r$ is the **dual norm** of $\|\cdot\|_q$ if $\frac{1}{q} + \frac{1}{r} = 1$ . Therefore, $r = 1 + \log(p)$ for $q = 1 + 1/\log(p)$ .



### Metrics

A metric on a set is a function that satisfies the minimal properties of a distance.

## Definition (Metric)

Let  $\mathcal{X}$  be a set, then a function  $d(\cdot, \cdot) : \mathcal{X} \times \mathcal{X} \to \mathbb{R}$  is a metric if  $\forall \mathbf{x}, \mathbf{y} \in \mathcal{X}$ :

- (a)  $d(\mathbf{x}, \mathbf{y}) \ge 0$  for all  $\mathbf{x}$  and  $\mathbf{y}$  (nonnegativity)
- (b)  $d(\mathbf{x}, \mathbf{y}) = 0$  if and only if  $\mathbf{x} = \mathbf{y}$  (definiteness)
- (c)  $d(\mathbf{x}, \mathbf{y}) = d(\mathbf{y}, \mathbf{x})$  (symmetry)
- (d)  $d(\mathbf{x}, \mathbf{y}) \le d(\mathbf{x}, \mathbf{z}) + d(\mathbf{z}, \mathbf{y})$  (triangle inequality)
- A pseudo-metric satisfies (a), (c) and (d) but not necessarily (b)
- A metric space  $(\mathcal{X}, d)$  is a set  $\mathcal{X}$  with a metric d defined on  $\mathcal{X}$
- Norms induce metrics while pseudo-norms induce pseudo-metrics

### Example

- Euclidean distance:  $d_E(\mathbf{x}, \mathbf{y}) = \|\mathbf{x} \mathbf{y}\|_2$
- Bregman distance:  $d_B(\cdot, \cdot)$  ...more on this later!





# **Basic matrix definitions**

# Definition (Nullspace of a matrix)

The nullspace of a matrix,  $\mathbf{A} \in \mathbb{R}^{n imes p}$ , (denoted by  $\mathrm{null}(\mathbf{A})$ ) is defined as

$$\operatorname{null}(\mathbf{A}) = \{\mathbf{x} \in \mathbb{R}^p \mid \mathbf{A}\mathbf{x} = \mathbf{0}\}$$

- null(A) is the set of vectors mapped to zero by A.
- $\operatorname{null}(\mathbf{A})$  is the set of vectors orthogonal to the rows of  $\mathbf{A}$ .

### Definition (Range of a matrix)

The range of a matrix,  $\mathbf{A} \in \mathbb{R}^{n \times p}$ , (denoted by  $\operatorname{range}(\mathbf{A})$ ) is defined as

range(
$$\mathbf{A}$$
) = { $\mathbf{A}\mathbf{x} \mid \mathbf{x} \in \mathbb{R}^p$ }  $\subseteq \mathbb{R}^n$ 

• range(A) is the span of the columns (or the column space) of A.

### Definition (Rank of a matrix)

The rank of a matrix,  $\mathbf{A} \in \mathbb{R}^{n \times p}$ , (denoted by  $\operatorname{rank}(\mathbf{A})$ ) is defined as

$$\operatorname{rank}(\mathbf{A}) = \operatorname{\mathbf{dim}}(\operatorname{range}(\mathbf{A}))$$

- ▶ rank(A) is the maximum number of independent columns (or rows) of A, ⇒ rank(A)  $\leq \min(n, p)$ .
- ▶  $\operatorname{rank}(\mathbf{A}) = \operatorname{rank}(\mathbf{A}^T)$ ; and  $\operatorname{rank}(\mathbf{A}) + \operatorname{dim}(\operatorname{null}(\mathbf{A})) = n$ .



## Matrix definitions contd.

# Definition (Eigenvalues & Eigenvectors)

The vector  $\mathbf{x}$  is an eigenvector of a square matrix  $\mathbf{A} \in \mathbb{R}^{n \times n}$  if  $\mathbf{A}\mathbf{x} = \lambda \mathbf{x}$  where  $\lambda \in \mathbb{R}$  is called an eigenvalue of  $\mathbf{A}$ .

• A scales its eigenvectors by it eigenvalues.

Definition (Singular values & singular vectors)

For  $\mathbf{A} \in \mathbb{R}^{n \times p}$  and *unit* vectors  $\mathbf{u} \in \mathbb{R}^n$  and  $\mathbf{v} \in \mathbb{R}^p$  if

$$\mathbf{A}\mathbf{v} = \sigma\mathbf{u}$$
 and  $\mathbf{A}^T\mathbf{u} = \sigma\mathbf{v}$ 

then  $\sigma \in \mathbb{R}$  ( $\sigma \ge 0$ ) is a singular value of  $\mathbf{A}$ ;  $\mathbf{v}$  and  $\mathbf{u}$  are the right singular vector and the left singular vector respectively of  $\mathbf{A}$ .

#### Definition (Symmetric matrix)

A matrix  $\mathbf{A} \in \mathbb{R}^{n \times n}$  is symmetric if  $\mathbf{A} = \mathbf{A}^T$ .

#### Lemma

The eigenvalues of a symmetric  $\mathbf{A}$  are real.

#### Proof.

Assume 
$$\mathbf{A}\mathbf{x} = \lambda \mathbf{x}$$
,  $\mathbf{x} \in \mathbb{C}^p$ ,  $\mathbf{x} \neq \mathbf{0}$ , then  $\overline{\mathbf{x}}^T \mathbf{A}\mathbf{x} = \overline{\mathbf{x}}^T (\mathbf{A}\mathbf{x}) = \overline{\mathbf{x}}^T (\lambda \mathbf{x}) = \lambda \sum_{i=1}^n |x_i|^2$   
but  $\overline{\mathbf{x}}^T \mathbf{A}\mathbf{x} = \overline{(\mathbf{A}\mathbf{x})}^T \mathbf{x} = \overline{\lambda} \sum_{i=1}^n |x_i|^2 \Rightarrow \lambda = \overline{\lambda}$  i.e.  $\lambda \in \mathbb{R}$ 



### Matrix definitions contd.

Definition (Positive semidefinite & positive definite matrices) A symmetric matrix  $\mathbf{A} \in \mathbb{R}^{n \times n}$  is positive semidefinite (denoted  $\mathbf{A} \succeq 0$ ) if  $\mathbf{x}^T \mathbf{A} \mathbf{x} \ge 0$  for all  $\mathbf{x} \neq \mathbf{0}$ ; while it is positive definite (denoted  $\mathbf{A} \succ 0$ ) if  $\mathbf{x}^T \mathbf{A} \mathbf{x} > 0$ 

- $\mathbf{A} \succeq 0$  iff all its eigenvalues are nonnegative i.e.  $\lambda_{\min}(\mathbf{A}) \ge 0$ .
- Similarly,  $\mathbf{A} \succ 0$  iff all its eigenvalues are **positive** i.e.  $\lambda_{\min}(\mathbf{A}) > 0$ .
- A is negative semidefinite if  $-A \succeq 0$ ; while A is negative definite if  $-A \succ 0$ .
- Semidefinite ordering of two symmetric matrices, A and B:  $A \succeq B$  if  $A B \succeq 0$ .

### Example (Matrix inequalities)

1. If 
$$\mathbf{A} \succeq 0$$
 and  $\mathbf{B} \succeq 0$ , then  $\mathbf{A} + \mathbf{B} \succeq 0$ 

- 2. If  $A \succeq B$  and  $C \succeq D$ , then  $A + C \succeq B + D$
- 3. If  $\mathbf{B} \leq 0$  then  $\mathbf{A} + \mathbf{B} \leq \mathbf{A}$
- 4. If  $\mathbf{A} \succeq 0$  and  $\alpha \ge 0$ , then  $\alpha \mathbf{A} \succeq 0$
- 5. If  $\mathbf{A} \succ 0$ , then  $\mathbf{A}^2 \succ 0$
- 6. If  $\mathbf{A} \succ 0$ , then  $\mathbf{A}^{-1} \succ 0$



## Matrix decompositions

## Definition (Eigenvalue decomposition)

The eigenvalue decomposition of a square matrix,  $\mathbf{A} \in \mathbb{R}^{n \times n}$ , is given by:

$$\mathbf{A} = \mathbf{X} \mathbf{\Lambda} \mathbf{X}^{-1}$$

- ▶ the columns of  $\mathbf{X} \in \mathbb{R}^{n imes n}$ , i.e.  $\mathbf{x}_i$ , are eigenvectors of  $\mathbf{A}$
- $\Lambda = \operatorname{diag}(\lambda_1, \lambda_2, \dots, \lambda_n)$  where  $\lambda_i$  (also denoted  $\lambda_i(\mathbf{A})$ ) are eigenvalues of  $\mathbf{A}$
- A matrix that admits this decomposition is therefore called diagonalizable matrix

#### Eigendecomposition of symmetric matrices

If  $\mathbf{A} \in \mathbb{R}^{n \times n}$  is symmetric, the decomposition becomes  $\mathbf{A} = \mathbf{U} \mathbf{\Lambda} \mathbf{U}^T$ where  $\mathbf{U} \in \mathbb{R}^{n \times n}$  is unitary (or orthonormal), i.e.  $\mathbf{U}^T \mathbf{U} = \mathbf{I}$  and  $\lambda_i$  are real

If we order  $\lambda_1 \geq \lambda_2 \geq \cdots \geq \lambda_n$ ,  $\lambda_i(\mathbf{A})$  becomes the  $i^{\text{th}}$  largest eigenvalue of  $\mathbf{A}$ :

- $\lambda_n(\mathbf{A}) = \lambda_{\min}(\mathbf{A})$  is the minimum eigenvalue of  $\mathbf{A}$
- $\lambda_1(\mathbf{A}) = \lambda_{\max}(\mathbf{A})$  is the maximum eigenvalue of  $\mathbf{A}$



## Matrix decompositions contd

# Definition (Determinant of a matrix)

The determinant of a square matrix  $\mathbf{A} \in \mathbb{R}^{p \times p}$ , denoted by  $det(\mathbf{A})$ , is given by:

$$\det(\mathbf{A}) = \prod_{i=1}^{p} \lambda_i$$

where  $\lambda_i$  are *eigenvalues* of **A**.





## Matrix decompositions contd

# Definition (Singular value decomposition)

The singular value decomposition (SVD) of a matrix,  $\mathbf{A} \in \mathbb{R}^{n \times p}$ , is given by:

$$\mathbf{A} = \mathbf{U} \mathbf{\Sigma} \mathbf{V}^T = \sum_{i=1}^r \sigma_i \mathbf{u}_i \mathbf{v}_i^T$$

- ▶  $\operatorname{rank}(\mathbf{A}) = r \leq \min(n, p)$  and  $\sigma_i$  is the  $i^{\mathsf{th}}$  singular value of  $\mathbf{A}$
- $\mathbf{u}_i$  and  $\mathbf{v}_i$  are the  $i^{\mathsf{th}}$  left and right singular vectors of  $\mathbf{A}$  respectively
- $\mathbf{U} \in \mathbb{R}^{n \times r}$  and  $\mathbf{V} \in \mathbb{R}^{p \times r}$  are unitary matrices (i.e.  $\mathbf{U}^T \mathbf{U} = \mathbf{I}$ )
- $\Sigma = \operatorname{diag}(\sigma_1, \sigma_2, \dots, \sigma_r)$  where  $\sigma_1 \ge \sigma_2 \ge \dots \ge \sigma_r \ge 0$
- $\begin{array}{l} \mathbf{v}_i \text{ are eigenvectors of } \mathbf{A}^T \mathbf{A}; \ \sigma_i = \sqrt{\lambda_i \left(\mathbf{A}^T \mathbf{A}\right)} \ (\text{and } \lambda_i \left(\mathbf{A}^T \mathbf{A}\right) = 0 \ \text{for } i > r) \\ \text{since} \quad \mathbf{A}^T \mathbf{A} = \left(\mathbf{U} \boldsymbol{\Sigma} \mathbf{V}^T\right)^T \left(\mathbf{U} \boldsymbol{\Sigma} \mathbf{V}^T\right) = \left(\mathbf{V} \boldsymbol{\Sigma}^2 \mathbf{V}^T\right) \end{array}$

• 
$$\mathbf{u}_i$$
 are eigenvectors of  $\mathbf{A}\mathbf{A}^T$ ;  $\sigma_i = \sqrt{\lambda_i (\mathbf{A}\mathbf{A}^T)}$  (and  $\lambda_i (\mathbf{A}\mathbf{A}^T) = 0$  for  $i > r$ )  
since  $\mathbf{A}\mathbf{A}^T = (\mathbf{U}\Sigma\mathbf{V}^T) (\mathbf{U}\Sigma\mathbf{V}^T)^T = (\mathbf{U}\Sigma^2\mathbf{U}^T)$ 



## Matrix decompositions contd

# Definition (LU)

The LU factorization of a nonsingular square matrix,  $\mathbf{A} \in \mathbb{R}^{p \times p}$ , is given by:

 $\mathbf{A}=\mathbf{PLU}$ 

where P is a permutation matrix<sup>1</sup>, L is lower triangular and U is upper triangular.

## Definition (QR)

The **QR** factorization of any matrix,  $\mathbf{A} \in \mathbb{R}^{n \times p}$ , is given by:

 $\mathbf{A}=\mathbf{Q}\mathbf{R}$ 

where  $\mathbf{Q} \in \mathbb{R}^{n \times n}$  is an orthonormal matrix, i.e.  $\mathbf{Q}^T \mathbf{Q} = \mathbf{I}$ , and  $\mathbf{R} \in \mathbb{R}^{n \times p}$  is upper triangular.

## Definition (Cholesky)

The **Cholesky factorization** of a positive definite and symmetric matrix,  $\mathbf{A} \in \mathbb{R}^{p \times p}$ , is given by:

$$\mathbf{A} = \mathbf{L}\mathbf{L}^T$$

where  $\mathbf{L}$  is a lower triangular matrix with positive entries on the diagonal.

 $^1$  A matrix  $\mathbf{P} \in \mathbb{R}^{p imes p}$  is **permutation** if it has only one 1 in each row and each column.



# **Complexity of matrix operations**

#### Complexity of matrix operations

The complexity or *cost* of an algorithm is expressed in terms of **floating-point operations** (flops) as a function of the *problem dimension*.

## Definition (floating-point operation)

A **floating-point operation** (flop) is one addition, subtraction, multiplication, or division of two floating-point numbers.





## Complexity of matrix operations contd

Table: Complexity examples: vector are in  $\mathbb{R}^p$ , matrices in  $\mathbb{R}^{n \times p}$  or  $\mathbb{R}^{p \times m}$  for square matrices

Operation	Complexity	Remarks		
vector addition	p flops			
vector inner product	2p-1 flops	or $pprox 2p$ for $p$ large		
matrix-vector product	n(2p-1) flops	or $\approx 2np$ for $p$ large		
		$2m$ if ${f A}$ is sparse with $m$ nonzeros		
matrix-matrix product	mn(2p-1) flops	or $\approx 2mnp$ for $p$ large		
		much less if ${f A}$ is sparse $^1$		
LU decomposition	$\frac{2}{3}p^3 + 2p^2$ flops	or $\frac{2}{3}p^3$ for $p$ large much less if ${f A}$ is sparse $^1$		
		much less if ${f A}$ is sparse $^1$		
Cholesky decomposition	$\frac{1}{3}p^3 + 2p^2$ flops	or $\frac{1}{3}p^3$ for p large		
		much less if ${f A}$ is sparse $^1$		
SVD	$C_1 n^2 p + C_2 p^3$ flops	$C_1 = 4$ , $C_2 = 22$ for R-SVD algo.		
Determinant	complexity of SVD			

<sup>1</sup> Complexity depends on p, no. of nonzeros in  $\mathbf{A}$  and the sparsity pattern.





## Matrix norms

Similar to vector norms, matrix norms are a metric over matrices:

#### Definition (Matrix norm)

 $\begin{array}{ll} \mathsf{A} \text{ norm of an } n \times p \text{ matrix is a map } \| \cdot \| : \mathbb{R}^{n \times p} \to \mathbb{R} \text{ such that for all matrices} \\ \mathbf{A}, \mathbf{B} \in \mathbb{R}^{n \times p} \text{ and scalar } \lambda \in \mathbb{R} \\ (a) & \|\mathbf{A}\| \geq 0 \text{ for all } \mathbf{A} \in \mathbb{R}^{n \times p} & \textit{nonnegativity} \\ (b) & \|\mathbf{A}\| = 0 \text{ if and only if } \mathbf{A} = \mathbf{0} & \textit{definitiveness} \\ (c) & \|\lambda\mathbf{A}\| = |\lambda| \|\mathbf{A}\| & \textit{homogeniety} \\ (d) & \|\mathbf{A} + \mathbf{B}\| \leq \|\mathbf{A}\| + \|\mathbf{B}\| & \textit{triangle inequality} \end{array}$ 

### Definition (Matrix inner product)

Matrix inner product is defined as follows

$$\langle \mathbf{A}, \mathbf{B} \rangle = \mathsf{trace} \left( \mathbf{A} \mathbf{B}^T \right)$$
.



### Linear operators

- Matrices are often given in an implicit form.
- It is convenient to think of them as linear operators.

### Proposition (Linear operators & matrices)

Any linear operator in finite dimensional spaces can be represented as a matrix.

### Example

Given matrices  $\mathbf{A}, \mathbf{B}$  and  $\mathbf{X}$  with compatible dimensions and the *linear operator*  $\mathcal{M} : \mathbb{R}^{n \times p} \to \mathbb{R}^{np}$ , a linear operator can define the following implicit mapping

$$\mathcal{M}(\mathbf{X}) \coloneqq \left(\mathbf{B}^T \otimes \mathbf{A}\right) \operatorname{vec}(\mathbf{X}) = \operatorname{vec}(\mathbf{A}\mathbf{X}\mathbf{B}),$$

where  $\otimes$  is the Kronecker product and  $\operatorname{vec} : \mathbb{R}^{n \times p} \to \mathbb{R}^{np}$  is yet another linear operator that vectorizes its entries. Note: Clearly, it is more efficient to compute  $\operatorname{vec}(\mathbf{AXB})$  than to perform the *matrix* 

multiplication  $(\mathbf{B}^T \otimes \mathbf{A})$  vec $(\mathbf{X})$ .



## Definition (Operator norm)

The operator norm between  $\ell_q$  and  $\ell_r$   $(1 \le q, r \le \infty)$  of a matrix  ${f A}$  is defined as

$$\|\mathbf{A}\|_{q \to r} = \sup_{\|\mathbf{x}\|_q \le 1} \|\mathbf{A}\mathbf{x}\|_r$$

### Problem

Show that  $\|\mathbf{A}\|_{2\to 2} = \|\mathbf{A}\|$  i.e.,  $\ell_2$  to  $\ell_2$  operator norm is the spectral norm.

## Solution



## Other examples

▶ The  $||A||_{\infty\to\infty}$  (norm induced by  $\ell_{\infty}$ -norm) also denoted  $||A||_{\infty}$ , is the max-row-sum norm:

$$\|\mathbf{A}\|_{\infty \to \infty} := \sup\{\|\mathbf{A}\mathbf{x}\|_{\infty} \mid \|\mathbf{x}\|_{\infty} \le 1\} = \max_{i=1,...,n} \sum_{j=1}^{p} |a_{ij}|.$$

▶ The  $\|A\|_{1 \to 1}$  (norm induced by  $\ell_1$ -norm) also denoted  $\|A\|_1$ , is the max-column-sum norm:

$$\|\mathbf{A}\|_{1 \to 1} := \sup\{\|\mathbf{A}\mathbf{x}\|_1 \mid \|\mathbf{x}\|_1 \le 1\} = \max_{i=1,\dots,p} \sum_{j=1}^n |a_{ij}|.$$



### Useful relation for operator norms

The following identity holds

$$\|\mathbf{A}\|_{q \to r} := \max_{\|\mathbf{z}\|_r \leq 1, \|\mathbf{x}\|_q = 1} \langle \mathbf{z}, \mathbf{A}\mathbf{x} \rangle = \max_{\|\mathbf{x}\|_{q'} \leq 1, \|\mathbf{z}\|_{r'} = 1} \langle \mathbf{A}^T \mathbf{z}, \mathbf{x} \rangle =: \|\mathbf{A}^T\|_{q' \to r'}$$

whenever 1/q + 1/q' = 1 = 1/r + 1/r'.

## Example

- 1.  $\|\mathbf{A}\|_{\infty \to 1} = \|\mathbf{A}^T\|_{1 \to \infty}$ .
- 2.  $\|\mathbf{A}\|_{2 \to 1} = \|\mathbf{A}^T\|_{2 \to \infty}$ .
- 3.  $\|\mathbf{A}\|_{\infty \to 2} = \|\mathbf{A}^T\|_{1 \to 2}$ .





### Computation of operator norms

- ▶ The computation of some operator norms is NP-hard\* [3]; these include:
  - 1.  $\|\mathbf{A}\|_{\infty \to 1}$
  - $2. \|\mathbf{A}\|_{2 \to 1}$
  - 3.  $\|\mathbf{A}\|_{\infty \to 2}$

#### But some of them are approximable [5]; these include

- 1.  $\|\mathbf{A}\|_{\infty \to 1}$  (via Gronthendieck factorization)
- 2.  $\|\mathbf{A}\|_{\infty \to 2}$  (via Pietzs factorization)
- \*: See Lecture 3.





• Similar to vector  $\ell_p$ -norms, we have Schatten q-norms for matrices.

# Definition (Schatten *q*-norms)

$$\|\mathbf{A}\|_q := \left(\sum_{i=1}^p (\sigma(\mathbf{A})_i)^q\right)^{1/q}$$
, where  $\sigma(\mathbf{A})_i$  is the  $i^{th}$  singular value of  $\mathbf{A}$ .

Example (with  $r = \min\{n, p\}$  and  $\sigma_i = \sigma(\mathbf{A})_i$ )

$$\begin{aligned} \|\mathbf{A}\|_{1} &= \|\mathbf{A}\|_{*} &:= \sum_{i=1}^{r} \sigma_{i} &\equiv \operatorname{trace}\left(\sqrt{\mathbf{A}^{T}\mathbf{A}}\right) \quad (\operatorname{Nuclear/trace}) \\ \|\mathbf{A}\|_{2} &= \|\mathbf{A}\|_{F} &:= \sqrt{\sum_{i=1}^{r} (\sigma_{i})^{2}} &\equiv \sqrt{\sum_{i=1}^{n} \sum_{j=1}^{p} |a_{ij}|^{2}} \quad (\operatorname{Frobenius}) \\ \|\mathbf{A}\|_{\infty} &= \|\mathbf{A}\| &:= \max_{i=1,\dots,r} \{\sigma_{i}\} &\equiv \max_{\mathbf{x}\neq \mathbf{0}} \frac{\|\mathbf{A}\mathbf{x}\|}{\|\mathbf{x}\|} \quad (\operatorname{Spectral/matrix}) \end{aligned}$$





## Problem (Rank-*r* approximation)

Find  $\underset{\mathbf{X}}{\operatorname{arg\,min}} \|\mathbf{X} - \mathbf{Y}\|_F$  subject to:  $\operatorname{rank}(\mathbf{X}) \leq r$ .





### Problem (Rank-r approximation)

 $\begin{array}{ll} \mathsf{Find} & \mathop{\arg\min}\limits_{\mathbf{X}} \ \|\mathbf{X}-\mathbf{Y}\|_F \ \ \mathsf{subject to:} \ \ \mathrm{rank}(\mathbf{X}) \leq r. \\ \end{array}$ 

## Solution (Eckart-Young-Mirsky Theorem)

$$\begin{aligned} \underset{\mathbf{X}:\operatorname{rank}(\mathbf{X})\leq r}{\arg\min} & \|\mathbf{X} - \mathbf{Y}\|_{F} = \underset{\mathbf{X}:\operatorname{rank}(\mathbf{X})\leq r}{\arg\min} & \|\mathbf{X} - \mathbf{U}\boldsymbol{\Sigma}_{\mathbf{Y}}\mathbf{V}^{T}\|_{F}, \quad (\mathsf{SVD}) \\ &= \underset{\mathbf{X}:\operatorname{rank}(\mathbf{X})\leq r}{\arg\min} & \|\mathbf{U}^{T}\mathbf{X}\mathbf{V} - \boldsymbol{\Sigma}_{\mathbf{Y}}\|_{F}, \quad (\mathsf{unit. invar. of } \|\cdot\|_{F}) \\ &= \mathbf{U}\left(\underset{\mathbf{X}:\operatorname{rank}(\mathbf{X})\leq r}{\arg\min} & \|\mathbf{X} - \boldsymbol{\Sigma}_{\mathbf{Y}}\|_{F}\right)\mathbf{V}^{T}, \quad (\mathsf{sparse approx.}) \\ &= \mathbf{U}H_{r}\left(\boldsymbol{\Sigma}_{\mathbf{Y}}\right)\mathbf{V}^{T}, \quad (r\text{-sparse approx. of the diagonal entries}) \end{aligned}$$

Singular value hard thresholding operator  $H_r$  performs the best rank-r approximation of a matrix via sparse approximation: We keep the r largest singular values of the matrix and set the rest to zero.





The last step of the above solution makes use of the Mirsky inequality.

# Theorem (Mirsky inequality)

If  $\mathbf{A}, \mathbf{B}$  are  $p \times p$  matrices with singular values

 $\sigma_1 \ge \sigma_2 \ge \cdots \ge \sigma_p \ge 0, \quad \tau_1 \ge \tau_2 \ge \cdots \ge \tau_p \ge 0$ 

respectively. Let  $\pmb{\sigma} = (\sigma_1, \ldots, \sigma_p)^T$  and  $\pmb{\tau} = (\tau_1, \ldots, \tau_p)^T$ , then

$$\|\mathbf{A} - \mathbf{B}\|_F \ge \|\boldsymbol{\sigma} - \boldsymbol{\tau}\|_2.$$

 Mirsky theorem is proved using the following simplified version of von Neumann trace inequality.

#### Theorem (von Neumann trace inequality)

If  $\mathbf{A}, \mathbf{B}$  are  $p \times p$  matrices with singular values

$$\sigma_1 \ge \sigma_2 \ge \cdots \ge \sigma_p \ge 0, \quad \tau_1 \ge \tau_2 \ge \cdots \ge \tau_p \ge 0$$

respectively. Let  $\pmb{\sigma} = (\sigma_1, \ldots, \sigma_p)^T$  and  $\pmb{\tau} = (\tau_1, \ldots, \tau_p)^T$ , then

$$\langle \mathbf{A}, \mathbf{B} 
angle \leq \langle oldsymbol{\sigma}, oldsymbol{ au} 
angle$$



Matrix & vec	tor norm and	alogy			
	Vectors	$\ \mathbf{x}\ _1$	$\ \mathbf{x}\ _2$	$\ \mathbf{x}\ _{\infty}$	
	Matrices	$\ \mathbf{X}\ _{*}$	$\ \mathbf{X}\ _F$	$\ \mathbf{X}\ $	

# Definition (Dual of a matrix)

The dual norm of  $\mathbf{A} \in \mathbb{R}^{n \times p}$  is defined as

$$\|\mathbf{A}\|^* = \sup \left\{ \operatorname{trace} \left( \mathbf{A}^T \mathbf{X} \right) \mid \|\mathbf{X}\| \leq 1 \right\}.$$

# Matrix & vector dual norm analogy

Vector primal norm	$\ \mathbf{x}\ _1$	$\ \mathbf{x}\ _2$	$\ \mathbf{x}\ _{\infty}$
Vector dual norm	$\ \mathbf{x}\ _{\infty}$	$\ \mathbf{x}\ _2$	$\ \mathbf{x}\ _1$
Matrix primal norm	$\ \mathbf{X}\ _{*}$	$\ \mathbf{X}\ _F$	X
Matrix dual norm	$\ \mathbf{X}\ $	$\ \mathbf{X}\ _F$	$\ \mathbf{X}\ _{*}$





# Definition (Nuclear norm computation)

 $\|\mathbf{A}\|_* := \|\boldsymbol{\sigma}(\mathbf{A})\|_1 \quad \text{where } \boldsymbol{\sigma}(\mathbf{A}) \text{ is a vector of singular values of } \mathbf{A}$ 

$$= \min_{\mathbf{U}, \mathbf{V}: \mathbf{A} = \mathbf{U}\mathbf{V}^H} \|\mathbf{U}\|_F \|\mathbf{V}\|_F = \min_{\mathbf{U}, \mathbf{V}: \mathbf{A} = \mathbf{U}\mathbf{V}^H} \frac{1}{2} \left( \|\mathbf{U}\|_F^2 + \|\mathbf{V}\|_F^2 \right)$$

Additional useful properties are below:

- ▶ Nuclear vs. Frobenius:  $\|\mathbf{A}\|_F \leq \|\mathbf{A}\|_* \leq \sqrt{\mathsf{rank}(\mathbf{A})} \cdot \|\mathbf{A}\|_F$
- Hölder for matrices:  $|\langle \mathbf{A}, \mathbf{B} \rangle| \le \|\mathbf{A}\|_p \|\mathbf{B}\|_q$ , when  $\frac{1}{p} + \frac{1}{q} = 1$
- We have

$$\begin{array}{ll} 1. & \|\mathbf{A}\|_{2 \rightarrow 2} \leq \|\mathbf{A}\|_{F} \\ 2. & \|\mathbf{A}\|_{2 \rightarrow 2}^{2} \leq \|\mathbf{A}\|_{1 \rightarrow 1} \|\mathbf{A}\|_{\infty \rightarrow \infty} \\ 3. & \|\mathbf{A}\|_{2 \rightarrow 2}^{2} \leq \|\mathbf{A}\|_{1 \rightarrow 1} \text{ when } \mathbf{A} \text{ is self-adjoint} \end{array}$$



## \*Matrix perturbation inequalities

▶ In the theorems below  $\mathbf{A}, \mathbf{B} \in \mathbb{R}^{p \times p}$  are symmetric matrices with spectra  $\{\lambda_i(\mathbf{A})\}_{i=1}^p$  and  $\{\lambda_i(\mathbf{B})\}_{i=1}^p$  where  $\lambda_1 \ge \lambda_2 \ge \cdots \ge \lambda_p$ .

Theorem (Lidskii inequality)

$$\begin{split} \lambda_{i_1} \left( \mathbf{A} + \mathbf{B} \right) + \cdots + \lambda_{i_n} \left( \mathbf{A} + \mathbf{B} \right) &\leq \lambda_{i_1} \left( \mathbf{A} \right) + \cdots \lambda_{i_n} \left( \mathbf{A} \right) + \lambda_{i_1} \left( \mathbf{B} \right) + \cdots + \lambda_{i_n} \left( \mathbf{B} \right), \\ \text{for any } 1 &\leq i_1 \leq \cdots \leq i_n \leq p. \end{split}$$

Theorem (Weyl inequality)

 $\lambda_{i+j-1}\left(\mathbf{A}+\mathbf{B}\right) \leq \lambda_{i}\left(\mathbf{A}\right) + \lambda_{j}\left(\mathbf{B}\right), \quad \text{for any } i,j \geq 1 \ \text{ and } i+j-1 \leq p.$ 

Theorem (Interlacing property)

Let  $\mathbf{A}_n = \mathbf{A}(1:n,1:n)$ , then  $\lambda_{n+1} (\mathbf{A}_{n+1}) \leq \lambda_n (\mathbf{A}_n) + \lambda_n (\mathbf{A}_{n+1})$  for  $n = 1, \dots, p$ .

- These inequalities **hold** in the more general setting when  $\lambda_i$  are replaced by  $\sigma_i$ .
- ► The list goes on to include Wedins bounds, Wielandt-Hoffman bounds and so on.
- More on such inequalities can be found in Terry Tao's blog (254A, Notes 3a).





## \*Tensors

Tensors provide a natural and concise mathematical represention of data.

# Definition (Tensor)

An  $m^{\text{th}}$ -rank tensor in p-dimensional space is a mathematical object that has p indices and  $p^m$  components and obeys certain transformation rules.

- ► In the literature, order is used interchangeably with rank, i.e.,  $k^{th}$ -rank tensor is also referred to as an order-k tensor.
- Tensors are multidimensional arrays and are a generalization of:
  - 1. scalars tensors with no indices; i.e., zeroth-rank tensor.
  - 2. vectors tensors with exactly one index; i.e., first-rank tensor.
  - 3. matrices tensors with exactly two indices; i.e., second-rank tensor.
- Think of the third-order Taylor series expansion



## \*Tensors contd.

### Caveat!

Not much is known about tensors and the generalizability of matrix notions to tensors:

- The notion of tensor (symmetric) rank is considerably more delicate than matrix (symmetric) rank. For instance:
  - 1. Not clear a priori that the symmetric rank should even be finite [2].
  - 2. Removal of the best rank-1 approximation of a general tensor may increase the tensor rank of the residual [4].
- It is NP-hard to compute the rank of a tensor in general; only approximations of (super) symmetric tensors possible [1].





### References

- Anima Anandkumar, Rong Ge, Daniel Hsu, Sham M Kakade, and Matus Telgarsky. Tensor decompositions for learning latent variable models. arXiv preprint arXiv:1210.7559, 2012.
- [2] Pierre Comon, Gene Golub, Lek-Heng Lim, and Bernard Mourrain. Symmetric tensors and symmetric tensor rank. SIAM Journal on Matrix Analysis and Applications, 30(3):1254–1279, 2008.
- [3] Simon Foucart and Holger Rauhut.
   A mathematical introduction to compressive sensing. Springer, 2013.
- [4] Alwin Stegeman and Pierre Comon.

Subtracting a best rank-1 approximation may increase tensor rank. Linear Algebra and its Applications, 433(7):1276–1300, 2010.

[5] Joel A Tropp.

Column subset selection, matrix factorization, and eigenvalue optimization.

In Proceedings of the Twentieth Annual ACM-SIAM Symposium on Discrete Algorithms, pages 978-986, 2009.



