

Mathematics of Data: From Theory to Computation

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Outline

► Today

1. Convex constrained optimization

- Problem setting, common structures and basis assumptions
- Solutions and approximate solutions
- Motivating examples

2. Optimality and duality

- Optimality condition
- Lagrange dualization
- Min-max formulation
- Equivalent interpretations of optimality condition.
- Dual decomposition ability

3. Classical solution methods

- Convex problem with equality constraints and null space method.
- Projected gradient method
- Frank-Wolfe method
- Quadratic penalty methods
- Augmented Lagrangian methods
- Alternating minimization algorithm (AMA)
- Alternating direction method of multipliers (ADMM)

4. Next week

1. Nonsmooth constrained optimization

Reading material

1. S. Boyd and L. Vandenberghe, "*Convex Optimization*", University Press, Cambridge, 2004.
 - ▶ Chapter 4 – Convex optimization problems
 - ▶ Chapter 5 – Duality
 - ▶ Section 10.1-Chapter 10 – Equality constrained minimization.
2. J. Nocedal and S. Wright, "*Numerical Optimization*", Springer-Verlag, 1999.
 - ▶ Chapter 17 – Penalty, Barrier and augmented Lagrangian methods, Section 17.4.
3. S. Boyd, N. Parikh, E. Chu, B. Peleato, and J. Eckstein, "*Distributed Optimization and Statistical Learning via the Alternating Direction Method of Multipliers*", Foundations and Trends in Machine Learning, 3(1):1–122, 2011.

Motivation

Motivation

- ▶ Unknown **parameters** in a **model** are **constrained** in practice.
- ▶ **Constrained convex optimization formulations** naturally encode these constraints.
- ▶ Hence, this lecture develops **numerical methods** for **constrained convex optimization**.

Mathematical form of constrained convex optimization

General setting of constrained convex optimization problems

$$f^* := \begin{cases} \min_{\mathbf{x} \in \mathbb{R}^p} & f(\mathbf{x}) \\ \text{s.t.} & \mathbf{A}\mathbf{x} = \mathbf{b}, \\ & \mathbf{x} \in \mathcal{X}. \end{cases} \quad (1)$$

- ▶ $f \in \mathcal{F}(\mathbb{R}^p)$ is a **convex** function
- ▶ $\mathbf{A} \in \mathbb{R}^{n \times p}$, $\mathbf{b} \in \mathbb{R}^n$
- ▶ \mathcal{X} is a nonempty, closed **convex** set.

Problem sources

- ▶ Many real-world **applications** (e.g., linear inverse problems, matrix completion) can be directly formulated as (1).
- ▶ Often times, computational considerations lead to (1) by **reformulations** of **existing unconstrained problems** (e.g., composite convex minimization, consensus optimization, and convex splitting).
- ▶ Many **standard convex optimization** formulations naturally fall under (1), such as *linear programming, convex quadratic programming, second order cone programming, semidefinite programming and geometric programming*.

Structures of constrained convex optimization

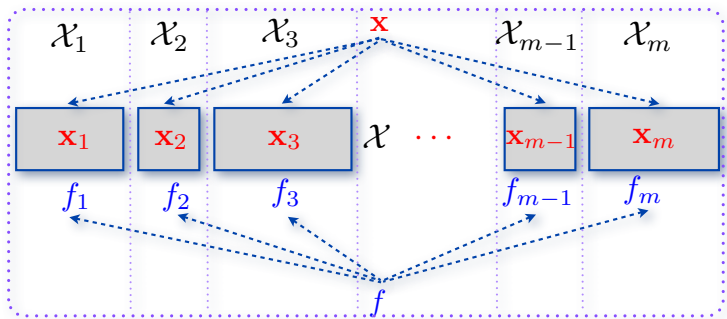
Common structures

When **designing** a **numerical solution method** for solving **problem (1)**, we must **rely** on individual **structures** of f and \mathcal{X} .

In this lecture, we mainly rely on the following two structures:

- ▶ **Decomposability** of f and \mathcal{X} .
- ▶ **Tractable proximity**

Decomposability illustration



Decomposability and tractable proximity

Decomposable structure

The function f and the feasible set \mathcal{X} have the following structure

$$f(\mathbf{x}) := \sum_{i=1}^m f_i(\mathbf{x}_i), \quad \text{and} \quad \mathcal{X} := \mathcal{X}_1 \times \cdots \times \mathcal{X}_m.$$

where $m \geq 1$ is the **number of components**, \mathbf{x}_i is a **sub-vector** (component) of \mathbf{x} , $f_i : \mathbb{R}^{p_i} \rightarrow \mathbb{R} \cup \{+\infty\}$ is **convex** and $\sum_{i=1}^m p_i = p$.

Decomposability and tractable proximity

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Tractable proximity

- ▶ Each **component** f_i has a **'tractable proximal operator'** ($i = 1, \dots, m$).
- ▶ The component **feasible set** \mathcal{X}_i has **simple projection** ("tractable proximity" of the **indicator function** of \mathcal{X}_i).

Solutions and solution set

Definition (Feasible set)

The set

$$\mathcal{D} := \{\mathbf{x} \in \mathbb{R}^p : \mathbf{x} \in \mathcal{X}, \mathbf{Ax} = \mathbf{b}\} \quad (2)$$

is called the **feasible set** of (1). Any point $\mathbf{x} \in \mathcal{D}$ is called a **feasible point**.

Note: It is important to exclude the following trivial and pathological cases:

- ▶ $\mathcal{D} = \emptyset$, which leads to no solution of (1).
- ▶ $\mathcal{D} = \{\hat{\mathbf{x}}\}$, which leads to the unique solution $\mathbf{x}^* = \hat{\mathbf{x}}$ of (1).

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Definition (Solution)

A feasible point $\mathbf{x}^* \in \mathcal{D}$ is called a **globally optimal solution** (or solution) of (1) if

$$f(\mathbf{x}^*) \leq f(\mathbf{x}), \quad \forall \mathbf{x} \in \mathcal{D}.$$

All **solutions** of (1) forms the **solution set** \mathcal{S}^* of (1).

Note:

- ▶ The solution set \mathcal{S}^* is **closed and convex**.
- ▶ If \mathbf{x} is **not feasible**, one may have $f(\mathbf{x}) \leq f^*$ in the **constrained setting case**.

Approximate solution

Solution certification

- ▶ Computing an **exact solution** $\mathbf{x}^* \in \mathcal{S}^*$ is **impracticable** unless problem has a **closed form solution** (which is very limited in reality).
- ▶ We can only compute a point \mathbf{x}_ϵ^* that **approximates** \mathbf{x}^* up to a given **accuracy** ϵ in a **given sense** by using **numerical optimization algorithms**.

There are **several ways** of certifying an **approximate solution**. We use the following definition.

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Definition (Approximate solution)

Given a **tolerance** $\epsilon \geq 0$, a point $\mathbf{x}_\epsilon^* \in \mathbb{R}^p$ is called an **ϵ -solution** of (1) if

$$\begin{cases} |f(\mathbf{x}_\epsilon^*) - f^*| \leq \epsilon & \text{(objective residual),} \\ \|\mathbf{A}\mathbf{x}_\epsilon^* - \mathbf{b}\| \leq \epsilon & \text{(feasibility gap),} \\ \mathbf{x}_\epsilon^* \in \mathcal{X} & \text{(exact feasibility).} \end{cases}$$

Very often, \mathcal{X} is a “**simple set**.” Hence, checking $\mathbf{x}_\epsilon^* \in \mathcal{X}$ is **acceptable** in practice.

Motivating example: Composite convex minimization

Composite convex minimization

With a slight change in notation, let us recall the **composite convex minimization** problem in Lecture 5:

$$F^* := \min_{\mathbf{u} \in \mathbb{R}^p} \{F(\mathbf{u}) := h(\mathbf{u}) + g(\mathbf{u})\}, \quad (3)$$

where both g and h are closed and **convex**.

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Optimization reformulation

By **duplicating the variable** $\mathbf{v} = \mathbf{u}$, we can reformulate (3) as

$$\begin{array}{ll} \min & \{f(\mathbf{x}) := h(\mathbf{v}) + g(\mathbf{u})\} \\ \mathbf{x} := [\mathbf{u}, \mathbf{v}] \in \mathbb{R}^{2p} & \\ \text{s.t.} & \mathbf{u} - \mathbf{v} = 0. \end{array} \quad (4)$$

This problem falls into the form (1) with **separable objective** function f and $\mathcal{X} = \mathbb{R}^{2p}$. **The methods** studied in this lecture can also be used to solve the **composite convex problem** (3).

Image denoising/deblurring

Problem (Imaging denoising/deblurring)

Given an *observed image* $\mathbf{b} \in \mathbb{R}^{n \times p}$, the aim is to recover the *clean image* \mathbf{u} via $\mathbf{b} = \mathcal{A}(\mathbf{u}) + \mathbf{w}$, where \mathcal{A} is a *linear operator* and \mathbf{w} is a *Gaussian noise*.

Optimization formulation

$$\min_{\mathbf{u} \in \mathbb{R}^{n \times p}} \left\{ (1/2) \|\mathcal{A}(\mathbf{u}) - \mathbf{b}\|_F^2 + \rho \|\mathbf{D}\mathbf{u}\|_1 \right\} \quad (5)$$

where $\rho > 0$ is a *regularization parameter* and \mathbf{D} is given matrix.
By reformulating (5) as

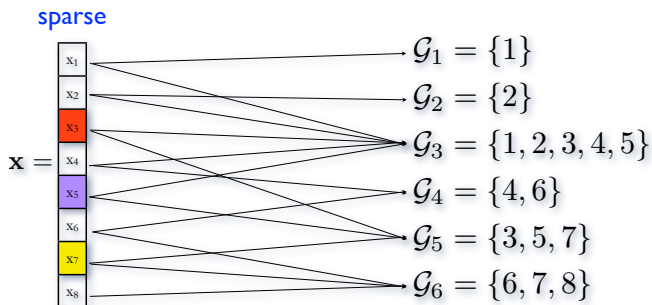
$$\begin{array}{ll} \min_{\mathbf{u} \in \mathbb{R}^{n \times p}} & \left\{ (1/2) \|\mathcal{A}(\mathbf{u}) - \mathbf{b}\|_F^2 + \rho \|\mathbf{v}\|_1 \right\} \\ \text{s.t.} & \mathbf{D}\mathbf{u} - \mathbf{v} = 0. \end{array} \quad (6)$$

This problem is of the form (1) with $\mathbf{x} := (\mathbf{u}^T, \mathbf{v}^T)^T$, $\mathcal{X} = \mathbb{R}^{np+nDp}$ and $f(\mathbf{x}) := (1/2) \|\mathcal{A}(\mathbf{u}) - \mathbf{b}\|_F^2 + \rho \|\mathbf{v}\|_1$.

Group sparse recovery

Sparse recovery

- ▶ Let $\mathcal{I} := \{1, \dots, p\}$ be the **set of indices**. Let $\mathcal{G} := \{\mathcal{G}_1, \dots, \mathcal{G}_m\}$ be the **set of m groups** $\mathcal{G}_i \subseteq \mathcal{I}$ and $\mathcal{I} \subseteq \cup_{i=1}^m \mathcal{U}_i$.
- ▶ For given **group** \mathcal{G}_i , and a vector $\mathbf{x} \in \mathbb{R}^p$, we use $\mathbf{x}_{\mathcal{G}_i} = \{x_j : j \in \mathcal{G}_i\}$.
- ▶ For fixed **group structure** \mathcal{G} , $\mathbf{x} \in \mathbb{R}^p$ is called **group sparse vector** if the **number of groups** in \mathcal{G} is **small**.
- ▶ Given a **linear operator** \mathbf{A} and an **observed/measurement** vector $\mathbf{b} \in \mathbb{R}^n$. We want to recover the **group sparse** input vector $\mathbf{x} \in \mathbb{R}^p$ such that $\mathbf{b} = \mathbf{A}\mathbf{x}$.



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Optimization formulation

$$\begin{aligned} \min_{\mathbf{x} \in \mathbb{R}^p} \quad & \sum_{\mathcal{G}_i \in \mathcal{G}} \|\mathbf{x}_{\mathcal{G}_i}\|_2 \\ \text{s.t.} \quad & \mathbf{A}\mathbf{x} = \mathbf{b}. \end{aligned} \tag{7}$$

Here, $f(\mathbf{x}) := \sum_{\mathcal{G}_i \in \mathcal{G}} \|\mathbf{x}_{\mathcal{G}_i}\|_2$ and $\mathcal{X} := \mathbb{R}^p$. This problem possesses two common structures: **decomposability** and **tractable proximity**.

When $m = p$ and $\mathcal{G}_i = \{i\}$, (7) reduces to the well-known **linear sparse recovery problem** (basis pursuit):

$$\min_{\mathbf{x} \in \mathbb{R}^p} \|\mathbf{x}\|_1 \quad \text{s.t.} \quad \mathbf{A}\mathbf{x} = \mathbf{b}. \tag{8}$$

Robust principle component analysis

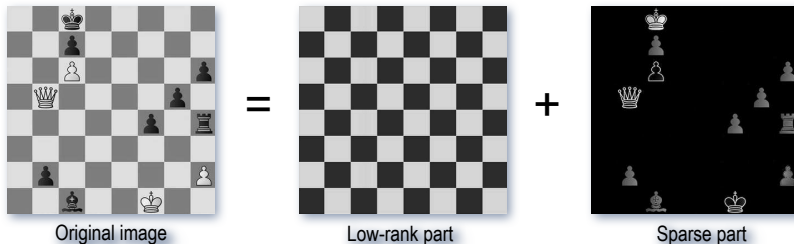
Robust principle component analysis (RPCA)

Assume that we are given a **large-scale input matrix** $\mathbf{M} \in \mathbb{R}^{m \times n}$, which can be **decomposed** as $\mathbf{M} = \mathbf{L}_0 + \mathbf{S}_0$, where \mathbf{L}_0 has **low-rank** and \mathbf{S}_0 is **sparse**. We **do not know** \mathbf{L}_0 and \mathbf{S}_0 and want to **recover** them given that they are **low-rank** and **sparse**, respectively.

Robust principle component analysis

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Motivating example: Robust principle component analysis

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Optimization formulation

$$\begin{aligned} \min_{\mathbf{L}, \mathbf{S} \in \mathbb{R}^{m \times n}} \quad & \|\text{vec}(\mathbf{S})\|_1 + \rho \|\mathbf{L}\|_*, \\ \text{s.t.} \quad & \mathbf{S} + \mathbf{L} = \mathbf{M}. \end{aligned} \quad (9)$$

Here $\rho > 0$ is a **weighted parameter** to trade-off between the **sparse** and **low-rank** terms, vec is the vectorization operator and $\|\cdot\|_*$ is the nuclear norm.

By letting

- ▶ $\mathbf{x} = [\mathbf{x}_1, \mathbf{x}_2] := [\text{vec}(\mathbf{S}), \text{vec}(\mathbf{L})]$
- ▶ $f(\mathbf{x}) = f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2) := \|\text{vec}(\mathbf{S})\|_1 + \rho \|\mathbf{L}\|_*$
- ▶ $\mathbf{A} = [\mathbf{I}, \mathbf{I}]$, $\mathbf{b} := \text{vec}(\mathbf{M})$ and
- ▶ $\mathcal{X} := \mathbb{R}^{mn}$.

Then, (9) can be transformed into (1).

Motivating example: Robust principle component analysis (cont)

Example - RPCA for object separation from video

Let \mathbf{M} be the matrix extracted from a video clip. Our aim is to **separate** objects (e.g., humans) and backgrounds by solving (9).

Motivating example: Robust principle component analysis (cont)

Example - RPCA for object separation from video

Let M be the matrix extracted from a video clip. Our aim is to **separate** objects (e.g., humans) and backgrounds by solving (9).

Result: One frame from the solution of (9)

One original image M



The low-rank part L



The sparse part S



Matrix completion

Matrix completion

Aim: Recover the **unknown entries** of a matrix $\mathbf{M} \in \mathbf{C}^{m \times n}$, when we only **observe** a **few** $q < m \times n$ entries at a **given locations** $(i, j) \in \Omega$.

Low-rankness: Since this is an **underdetermined problem**, there exist many matrix \mathbf{X} such that $\mathbf{X}_{ij} = \mathbf{M}_{ij}$ for all $(i, j) \in \Omega$. We would like to recover a **low-rank matrix** \mathbf{X} such that $\mathbf{X}_{ij} = \mathbf{M}_{ij}$ for all $(i, j) \in \Omega$.

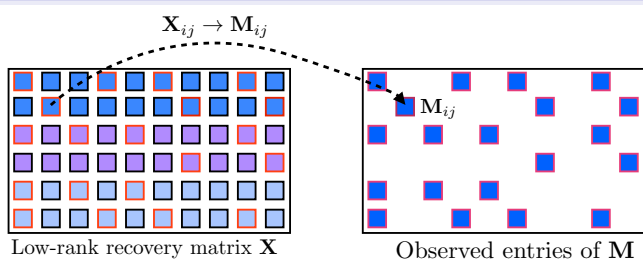
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Illustration



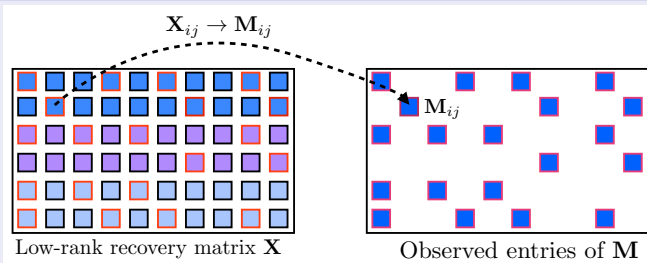
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Illustration



Convex relaxation of matrix completion

$$\begin{aligned}
 & \min_{\mathbf{X} \in \mathbb{C}^{m \times n}} && \|\mathbf{X}\|_* \\
 & \text{s.t.} && \mathbf{X}_{ij} = \mathbf{M}_{ij}, \forall (i, j) \in \Omega.
 \end{aligned} \tag{10}$$

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1. Nonsmooth constrained optimization

Optimality condition

Lagrange function

$$\mathcal{L}(\mathbf{x}, \lambda) := f(\mathbf{x}) + \lambda^T (\mathbf{A}\mathbf{x} - \mathbf{b}).$$

Here, $\lambda \in \mathbb{R}^n$ is the vector of **Lagrange multipliers** (or **dual** variables) w.r.t. $\mathbf{A}\mathbf{x} = \mathbf{b}$.

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Optimality condition

The **optimality condition** of (1) can be written as

$$\begin{cases} 0 & \in \mathbf{A}^T \lambda^* + \partial f(\mathbf{x}^*) + \mathcal{N}_{\mathcal{X}}(\mathbf{x}^*), \\ 0 & = \mathbf{A}\mathbf{x}^* - \mathbf{b}. \end{cases} \quad (11)$$

Here:

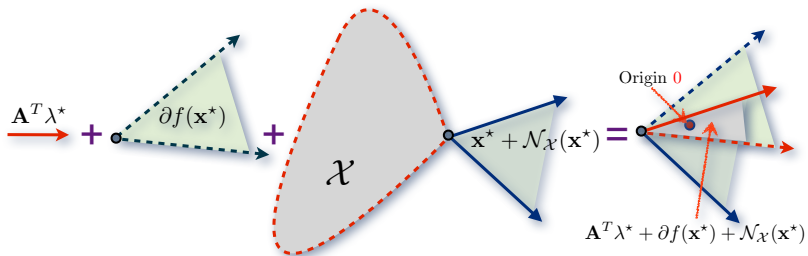
- ▶ $\partial f(\mathbf{x}) := \{\mathbf{z} \in \mathbb{R}^p : f(\mathbf{y}) \geq f(\mathbf{x}) + \mathbf{z}^T(\mathbf{y} - \mathbf{x}), \forall \mathbf{y} \in \mathbb{R}^p\}$ is the **subdifferential** of f at \mathbf{x} (see Lecture 2).
- ▶ $\mathcal{N}_{\mathcal{X}}$ is the normal cone of \mathcal{X} at \mathbf{x} defined as

$$\mathcal{N}_{\mathcal{X}}(\mathbf{x}) := \begin{cases} \{\mathbf{z} \in \mathbb{R}^p : \mathbf{z}^T(\mathbf{x} - \mathbf{y}) \geq 0, \forall \mathbf{y} \in \mathcal{X}\} & \text{if } \mathbf{x} \in \mathcal{X}, \\ \emptyset, & \text{if } \mathbf{x} \notin \mathcal{X}. \end{cases}$$

The condition (11) can be considered as the **KKT** (Karush-Kuhn-Tucker) condition. Any point $(\mathbf{x}^*, \lambda^*)$ satisfying (11) is called a **KKT point**. \mathbf{x}^* is called a **stationary point** and λ^* is the corresponding **multipliers**.

Example: Illustration

- ▶ This figure illustrates the first condition $0 \in \mathbf{A}^T \lambda^* + \partial f(\mathbf{x}^*) + \mathcal{N}_{\mathcal{X}}(\mathbf{x}^*)$.



Example: Basis pursuit

Example (Basis pursuit)

$$\min_{\mathbf{x} \in \mathbb{R}^p} \|\mathbf{x}\|_1 \quad \text{s.t. } \mathbf{A}\mathbf{x} = \mathbf{b}.$$

Note:

- ▶ $f(\mathbf{x}) := \|\mathbf{x}\|_1$ is **nonsmooth**, for any $\mathbf{v} \in \partial f(\mathbf{x})$ we have $v_i = +1$ if $x_i > 0$, $v_i = -1$ if $x_i < 0$ and $v_i \in (-1, 1)$ if $x_i = 0$.
- ▶ Since $\mathcal{X} \equiv \mathbb{R}^p$, we have $\mathcal{N}_{\mathcal{X}}(\mathbf{x}) = \{0\}$ for all \mathbf{x} .

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Optimality condition

The **optimality condition** of (11) becomes

$$\begin{cases} 0 \in \partial f(\mathbf{x}^*) + \mathbf{A}^T \lambda^* \\ 0 = \mathbf{A}\mathbf{x}^* - \mathbf{b}. \end{cases} \Leftrightarrow \begin{cases} (\mathbf{A}^T \lambda^*)_i = -1 & \text{if } x_i^* > 0, 1 \leq i \leq p \\ (\mathbf{A}^T \lambda^*)_i = +1 & \text{if } x_i^* < 0, 1 \leq i \leq p \\ (\mathbf{A}^T \lambda^*)_i \in (-1, 1) & \text{if } x_i^* = 0, 1 \leq i \leq p \\ \mathbf{A}\mathbf{x}^* = \mathbf{b}. \end{cases}$$

Min-max formulation and dual problem

Dual function and Dual problem

- **Dual function:**

$$d(\lambda) := \min_{\mathbf{x} \in \mathcal{X}} \{ \mathcal{L}(\mathbf{x}, \lambda) := f(\mathbf{x}) + \lambda^T (\mathbf{A}\mathbf{x} - \mathbf{b}) \}. \quad (12)$$

Let $\mathbf{x}^*(\lambda)$ be a **solution** of (12) then $d(\lambda)$ is finite if $\mathbf{x}^*(\lambda)$ **exists**. $d(\cdot)$ is concave and possibly nonsmooth.

- **Dual problem:** The following dual problem is **convex**

$$d^* := \max_{\lambda \in \mathbb{R}^n} d(\lambda) \quad (13)$$

Min-max formulation and dual problem

Dual function and Dual problem

- **Dual function:**

$$d(\lambda) := \min_{\mathbf{x} \in \mathcal{X}} \{\mathcal{L}(\mathbf{x}, \lambda) := f(\mathbf{x}) + \lambda^T (\mathbf{A}\mathbf{x} - \mathbf{b})\}. \quad (12)$$

Let $\mathbf{x}^*(\lambda)$ be a **solution** of (12) then $d(\lambda)$ is finite if $\mathbf{x}^*(\lambda)$ **exists**. $d(\cdot)$ is concave and possibly nonsmooth.

- **Dual problem:** The following dual problem is **convex**

$$d^* := \max_{\lambda \in \mathbb{R}^n} d(\lambda) \quad (13)$$

Min-max formulation

$$\begin{aligned} d^* &= \max_{\lambda \in \mathbb{R}^n} d(\lambda) = \max_{\lambda \in \mathbb{R}^n} \min_{\mathbf{x} \in \mathcal{X}} \{f(\mathbf{x}) + \lambda^T (\mathbf{A}\mathbf{x} - \mathbf{b})\} \\ &\leq \min_{\mathbf{x} \in \mathcal{X}} \max_{\lambda \in \mathbb{R}^n} \{f(\mathbf{x}) + \lambda^T (\mathbf{A}\mathbf{x} - \mathbf{b})\} = \begin{cases} \min_{\mathbf{x} \in \mathcal{X}} f(\mathbf{x}) & \text{if } \mathbf{A}\mathbf{x} = \mathbf{b}, \\ +\infty & \text{otherwise} \end{cases} \end{aligned} \quad (14)$$

Here, the inequality is due to the **max-min theorem** [6].

Example: Strictly convex quadratic programming

Strictly convex quadratic programming

$$\begin{array}{ll} \min_{\mathbf{x} \in \mathbb{R}^p} & (1/2)\mathbf{x}^T \mathbf{H}\mathbf{x} + \mathbf{h}^T \mathbf{x} \\ \text{s.t.} & \mathbf{A}\mathbf{x} = \mathbf{b}. \end{array}$$

where \mathbf{H} is symmetric positive definite.

Example: Strictly convex quadratic programming

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$$\begin{aligned} \min_{\mathbf{x} \in \mathbb{R}^p} \quad & (1/2)\mathbf{x}^T \mathbf{H} \mathbf{x} + \mathbf{h}^T \mathbf{x} \\ \text{s.t.} \quad & \mathbf{A} \mathbf{x} = \mathbf{b}. \end{aligned}$$

where \mathbf{H} is symmetric positive definite.

Dual problem is also a strictly convex quadratic program

- ▶ Lagrange function $\mathcal{L}(\mathbf{x}, \lambda) := (1/2)\mathbf{x}^T \mathbf{H} \mathbf{x} + (\mathbf{A}^T \lambda + \mathbf{h})^T \mathbf{x} - \mathbf{b}^T \lambda$.
- ▶ Dual function:

$$d(\lambda) = \min_{\mathbf{x} \in \mathbb{R}^p} \{ (1/2)\mathbf{x}^T \mathbf{H} \mathbf{x} + (\mathbf{A}^T \lambda + \mathbf{h})^T \mathbf{x} - \mathbf{b}^T \lambda \}$$

- ▶ Since $\mathbf{x}^*(\lambda) = -\mathbf{H}^{-1}(\mathbf{A}^T \lambda + \mathbf{h})$, we can obtain $d(\lambda)$ explicitly as

$$d(\lambda) = -(1/2)\lambda^T (\mathbf{A} \mathbf{H}^{-1} \mathbf{A}^T) \lambda - (\mathbf{b} + \mathbf{A} \mathbf{H}^{-1} \mathbf{h})^T \lambda.$$

- ▶ Dual problem (unconstrained):

$$d^* := \max_{\lambda \in \mathbb{R}^n} d(\lambda) \Leftrightarrow \min_{\lambda \in \mathbb{R}^n} \frac{1}{2} \lambda^T (\mathbf{A} \mathbf{H}^{-1} \mathbf{A}^T) \lambda + (\mathbf{b} + \mathbf{A} \mathbf{H}^{-1} \mathbf{h})^T \lambda.$$

Example: Nonsmoothness of the dual function

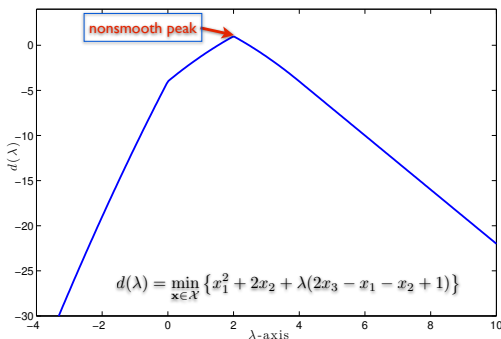
Consider a constrained convex problem:

$$\begin{aligned} \min_{\mathbf{x} \in \mathbb{R}^3} \quad & \{f(\mathbf{x}) := x_1^2 + 2x_2\}, \\ \text{s.t.} \quad & 2x_3 - x_1 - x_2 = 1, \\ & \mathbf{x} \in \mathcal{X} := [-2, 2] \times [-2, 2] \times [0, 2]. \end{aligned}$$

The **dual function** is defined as

$$d(\lambda) := \min_{\mathbf{x} \in \mathcal{X}} \{x_1^2 + 2x_2 + \lambda(2x_3 - x_1 - x_2 + 1)\}$$

is **concave** and **nonsmooth** as illustrated in the figure below.



Saddle point

Definition (Saddle point)

A point $(\mathbf{x}^*, \lambda^*) \in \mathcal{X} \times \mathbb{R}^n$ is called a **saddle point** of the Lagrange function \mathcal{L} if

$$\mathcal{L}(\mathbf{x}^*, \lambda) \leq \mathcal{L}(\mathbf{x}^*, \lambda^*) \leq \mathcal{L}(\mathbf{x}, \lambda^*), \quad \forall \mathbf{x} \in \mathcal{X}, \lambda \in \mathbb{R}^n.$$

Recall the minmax form:

$$\max_{\lambda} \min_{\mathbf{x} \in \mathcal{X}} \{ \mathcal{L}(\mathbf{x}, \lambda) := f(\mathbf{x}) + \lambda^T (\mathbf{A}\mathbf{x} - \mathbf{b}) \}. \quad ((12))$$

Saddle point

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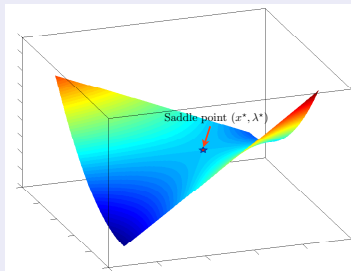
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Recall the minmax form:

$$\max_{\lambda} \min_{\mathbf{x} \in \mathcal{X}} \{ \mathcal{L}(\mathbf{x}, \lambda) := f(\mathbf{x}) + \lambda^T (\mathbf{Ax} - \mathbf{b}) \}. \quad ((12))$$

Illustration of saddle point: $\mathcal{L}(x, \lambda) := (1/2)x^2 + \lambda(x - 1)$ in \mathbb{R}^2



Slater's qualification condition

Slater's qualification condition

Recall $\text{relint}(\mathcal{X})$ the **relative interior** of the **feasible set** \mathcal{X} . The **Slater condition** requires

$$\text{relint}(\mathcal{X}) \cap \{\mathbf{x} : \mathbf{A}\mathbf{x} = \mathbf{b}\} \neq \emptyset. \quad (15)$$

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Recall $\text{relint}(\mathcal{X})$ the **relative interior** of the **feasible set** \mathcal{X} . The **Slater condition** requires

$$\text{relint}(\mathcal{X}) \cap \{\mathbf{x} : \mathbf{Ax} = \mathbf{b}\} \neq \emptyset. \quad (15)$$

Special cases

- ▶ If \mathcal{X} is **absent**, then (15) $\Leftrightarrow \boxed{\exists \bar{\mathbf{x}} : \mathbf{A}\bar{\mathbf{x}} = \mathbf{b}}$.
- ▶ If $\mathbf{Ax} = \mathbf{b}$ is **absent**, then (15) $\Leftrightarrow \boxed{\text{relint}(\mathcal{X}) \neq \emptyset}$.
- ▶ If $\mathbf{Ax} = \mathbf{b}$ is **absent** and $\mathcal{X} := \{\mathbf{x} : h(\mathbf{x}) \leq 0\}$, where h is $\mathbb{R}^p \rightarrow \mathbb{R}^q$ is convex, then

$$(15) \Leftrightarrow \boxed{\exists \bar{\mathbf{x}} : h(\bar{\mathbf{x}}) < 0.}$$

Example: Slater's condition

Example

Let us consider the feasible set $\mathcal{D}_\alpha := \mathcal{X} \cap \mathcal{A}_\alpha$ as

$$\mathcal{X} := \{\mathbf{x} \in \mathbb{R}^2 : x_1^2 + x_2^2 \leq 1\} \quad \mathcal{A}_\alpha := \{\mathbf{x} \in \mathbb{R}^2 : x_1 + x_2 = \alpha\},$$

where $\alpha \in \mathbb{R}$.

Example: Slater's condition

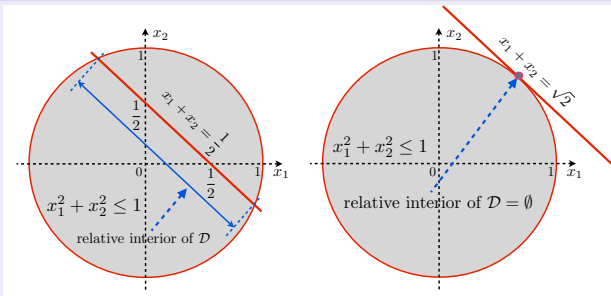
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Slater's condition holds and does not hold



$\mathcal{D}_{1/2}$ satisfies Slater's condition – $\mathcal{D}_{\sqrt{2}}$ -does not satisfy Slater's condition

Necessary and sufficient condition

Theorem (Necessary and sufficient optimality condition)

Under *Slater's condition* (15): $\text{relint}(\mathcal{X}) \cap \{\mathbf{x} : \mathbf{Ax} = \mathbf{b}\} \neq \emptyset$, the *KKT condition* (11)

$$\begin{cases} 0 \in \mathbf{A}^T \lambda^* + \partial f(\mathbf{x}^*) + \mathcal{N}_{\mathcal{X}}(\mathbf{x}^*), \\ 0 = \mathbf{Ax}^* - \mathbf{b}. \end{cases}$$

is *necessary and sufficient* for a point $(\mathbf{x}^*, \lambda^*) \in \mathcal{X} \times \mathbb{R}^n$ being an *optimal solution* for the primal problem (1) and dual problem (13):

$$f^* := \begin{cases} \min_{\mathbf{x} \in \mathbb{R}^p} & f(\mathbf{x}) \\ \text{s.t.} & \mathbf{Ax} = \mathbf{b}, \mathbf{x} \in \mathcal{X}, \end{cases} \quad \text{and} \quad d^* := \max_{\lambda \in \mathbb{R}^n} d(\lambda).$$

Necessary and sufficient condition

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Strong duality

- ▶ By definition of f^* and d^* , we always have $d^* \leq f^*$ (**weak duality**).
- ▶ Under Slater's condition and $\mathcal{X}^* \neq \emptyset$, we have $d^* = f^*$ (**strong duality**).
- ▶ Any solution $(\mathbf{x}^*, \lambda^*)$ of the KKT condition (11) is also a **saddle point**.

What happens if Slater's condition does not hold?

Without Slater's condition, KKT condition is only sufficient but not necessary, i.e., if $(\mathbf{x}^*, \lambda^*)$ satisfies the KKT condition, then \mathbf{x}^* is a global solution of (1) but not vice versa.

Example (Violating Slater's condition)

Consider the following constrained convex problem:

$$\min_{\mathbf{x} \in \mathbb{R}^2} \{x_1 : x_2 = 0, x_1^2 - x_2 \leq 0\}$$

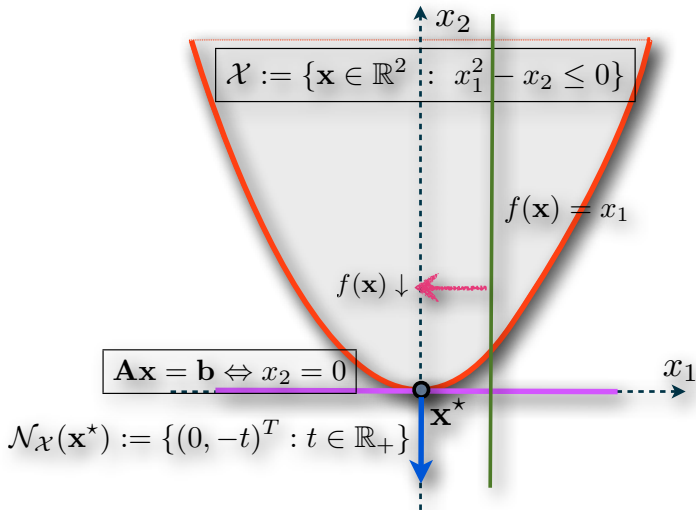
In the setting (1), we have $\mathbf{A} := [0, 1]$, $\mathbf{b} = 0$, $\mathcal{X} = \{\mathbf{x} \in \mathbb{R}^2 : x_1^2 - x_2 \leq 0\}$. The feasible set $\mathcal{D} := \{\mathbf{x} \in \mathbb{R}^2 : x_2 = 0, x_1^2 - x_2 \leq 0\} = \{(0, 0)^T\}$ contains only one point, which is also the optimal solution of the problem, i.e., $\mathbf{x}^* := (0, 0)^T$.

In this case, Slater's condition is definitely violated. Let us check the KKT condition. Since $\mathcal{N}_{\mathcal{X}}(\mathbf{x}^*) = \{(0, -t)^T : t \geq 0\}$, we can write the KKT condition as

$$\begin{bmatrix} 1 \\ 0 \end{bmatrix} + \begin{bmatrix} 0 \\ 1 \end{bmatrix} \lambda + \begin{bmatrix} 0 \\ -t \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}, \quad \lambda \in \mathbb{R}, t \in \mathbb{R}_+.$$

Since this linear system has no solution due to the first equation $1 = 0$, the KKT condition is inconsistent.

Violating Slater's condition



Variational inequality (VI) formulation

Primal-dual mapping

For simplicity, we assume that f is **smooth**. We introduce $\mathbf{z} := (\mathbf{x}^T, \lambda^T)^T \in \mathbb{R}^{p+n}$ and two mappings:

$$M(\mathbf{z}) := \begin{bmatrix} \nabla f(\mathbf{x}) + \mathbf{A}^T \lambda \\ \mathbf{A}\mathbf{x} - \mathbf{b} \end{bmatrix} \quad \text{and} \quad \mathcal{T}(\mathbf{z}) := \mathcal{N}_{\mathcal{X}}(\mathbf{x}) \times \{0^n\}. \quad (16)$$

Then $M : \mathbb{R}^{p+n} \rightarrow \mathbb{R}^{p+n}$ is a **single-valued mapping** and $\mathcal{T} : \mathbb{R}^{p+n} \rightrightarrows \mathbb{R}^{p+n}$ is a **set-valued mapping**.

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Inclusion and VI formulation

- ▶ The **optimality condition** (11) can be written as an **inclusion**:

$$0 \in \mathcal{R}(\mathbf{z}) := M(\mathbf{z}) + \mathcal{T}(\mathbf{z}).$$

- ▶ (11) can also be expressed as a **variational inequality**:

$$M(\mathbf{z}^*)^T (\mathbf{z} - \mathbf{z}^*) \geq 0, \quad \forall \mathbf{z} \in \mathcal{Z} := \mathcal{X} \times \mathbb{R}^n. \quad (17)$$

Dual decomposition ability

Roles of strong duality

- ▶ **Strong duality** is a **key property** in convex optimization, which creates a connection between **primal** problem (1) and **dual** problem (13).
- ▶ Under **Slater's condition**, **strong duality** holds, i.e., $f^* = d^*$.
- ▶ Principally, by solving **dual** problem (13), we can recover a **solution** of **primal** problem (1) and vice versa.

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Decomposability is a key property for parallel algorithms

- ▶ Under the **decomposable assumption**, the dual function d can be decomposed as

$$d(\lambda) = \sum_{i=1}^g d_i(\lambda) - \mathbf{b}^T \lambda.$$

where

$$d_i(\lambda) = \min_{\mathbf{x}_i \in \mathcal{X}_i} \{f_i(\mathbf{x}_i) + \lambda^T \mathbf{A}_i \mathbf{x}_i\}, \quad i = 1, \dots, g.$$

- ▶ Evaluating function $d_i(\cdot)$ and its [sub]gradients can be computed in **parallel**

Outline

▶ Today

1. Convex constrained optimization

- ▶ Problem setting, common structures and basis assumptions
- ▶ Solutions and approximate solutions
- ▶ Motivating examples

2. Optimality and duality

- ▶ Optimality condition
- ▶ Lagrange dualization
- ▶ Min-max formulation
- ▶ Equivalent interpretations of optimality condition.
- ▶ Dual decomposition ability

3. Classical solution methods

- ▶ Convex problem with equality constraints and null space method.
- ▶ Projected gradient method
- ▶ Frank-Wolfe method
- ▶ Quadratic penalty methods
- ▶ Augmented Lagrangian methods
- ▶ Alternating minimization algorithm (AMA)
- ▶ Alternating direction method of multipliers (ADMM)

4. Next week

1. Nonsmooth constrained optimization

Null space method for convex programs with equality constraints

Convex problems with equality constraints

We consider the case $\mathcal{X} \equiv \mathbf{R}^p$. Then (1) reduces to

$$f^* := \begin{cases} \min_{\mathbf{x} \in \mathbf{R}^p} & f(\mathbf{x}) \\ \text{s.t.} & \mathbf{Ax} = \mathbf{b}. \end{cases} \quad (18)$$

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Dimensional reduction

- ▶ Assume that $\text{rank}(\mathbf{A}) = m < p$, then the **dimension** of the **null space** $\dim(\text{null}(\mathbf{A})) = p - m$.
- ▶ By **eliminating** the equality constraints $\mathbf{A}\mathbf{x} = \mathbf{b}$, we can **reduce** the problem dimension from p to $p - m$.
- ▶ This elimination can be done via **projection** onto the **null space** $\text{null}(\mathbf{A})$ of \mathbf{A} , (e.g., by QR factorization of \mathbf{A}).
- ▶ Problem (18) can be transformed into an **unconstrained problem** with **dimension** $p - m$.

Null space method

Null space representation of the equality constraint $\mathbf{Ax} = \mathbf{b}$

- Any vector $\mathbf{x} \in \mathbb{R}^p$ can be represented as

$$\mathbf{x} = \bar{\mathbf{x}} + \mathbf{x}_{\mathcal{N}} = \bar{\mathbf{x}} + \mathbf{U}\mathbf{z},$$

where $\mathbf{x}_{\mathcal{N}} \in \text{null}(\mathbf{A})$, \mathbf{U} is a **basis** of $\text{null}(\mathbf{A})$ and $\bar{\mathbf{x}}$ satisfies $\mathbf{A}\bar{\mathbf{x}} = \mathbf{b}$.

- For any **feasible point** $\bar{\mathbf{x}}$ (i.e., $\mathbf{A}\bar{\mathbf{x}} = \mathbf{b}$), the point $\mathbf{x} := \bar{\mathbf{x}} + \mathbf{U}\mathbf{z}$ is also **feasible** to $\mathbf{Ax} = \mathbf{b}$, since

$$\mathbf{Ax} = \mathbf{A}\bar{\mathbf{x}} + \mathbf{A}\mathbf{U}\mathbf{z} = \mathbf{A}\bar{\mathbf{x}} = \mathbf{b}, \text{ since } \mathbf{A}\mathbf{U} = \mathbf{0}.$$

- \mathbf{U} can be computed via the **QR-factorization** of \mathbf{A}^T , and $\bar{\mathbf{x}}$ can be obtained by solving a **triangular linear system**.

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- \mathbf{U} can be computed via the **QR-factorization** of \mathbf{A}^T , and $\bar{\mathbf{x}}$ can be obtained by solving a **triangular linear system**.

Unconstrained formulation

By using the **null space representation** $\mathbf{x} = \bar{\mathbf{x}} + \mathbf{U}\mathbf{z}$, (18) can be transformed into the following **unconstrained formulation**:

$$\min_{\mathbf{z} \in \mathbb{R}^{p-n}} \left\{ \tilde{f}(\mathbf{z}) := f(\bar{\mathbf{x}} + \mathbf{U}\mathbf{z}) \right\}.$$

Example of null space representation

Problem

Given $s \in \mathbb{R}^3$, we want to compute the **projection** of s onto an **affine space** as:

$$\min_{\mathbf{x} \in \mathbb{R}^3} (1/2) \|\mathbf{x} - \mathbf{s}\|_2^2 \quad \text{s.t.} \quad \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & -1 \end{bmatrix} \mathbf{x} = \begin{bmatrix} 2 \\ 1 \end{bmatrix}, \quad \mathbf{x} \in \mathbb{R}^3. \quad (19)$$

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Null-space representation

- ▶ By computing the **QR factorization** of \mathbf{A}^T we obtain a 3×3 **orthonormal matrix** \mathbf{Z} and a 1×1 **triangular matrix** \mathbf{R} .
- ▶ Since $\text{rank}(\mathbf{A}) = 2$, $\dim(\text{null}(\mathbf{A})) = 3 - 2 = 1$, we take **the last column** of \mathbf{Z} to form a basis \mathbf{U} of $\text{null}(\mathbf{A})$, which is $\mathbf{U} := \begin{bmatrix} -\sqrt{2}/2 \\ \sqrt{2}/2 \\ 0 \end{bmatrix}$.
- ▶ The **two first columns** of \mathbf{Z} forms the basis of the **range space** of \mathbf{A}^T called \mathbf{V} .
- ▶ By solving $\mathbf{R}^T \mathbf{y} = \mathbf{b}$ we obtain $\mathbf{y} \approx (-1.15470, -0.20412)^T$. Therefore

$$\bar{\mathbf{x}} := \mathbf{V}\mathbf{y} = (3/4, 3/4, 1/2)^T.$$

- ▶ We finally obtain $\mathbf{x} = \bar{\mathbf{x}} + \mathbf{U}\mathbf{z}$, where $\mathbf{z} \in \mathbb{R}^2$ such that $\mathbf{A}\mathbf{x} = \mathbf{b}$.

From constrained to unconstrained formulation

The projection of \mathbf{s} onto the affine space $\mathbf{Ax} = \mathbf{b}$

Problem (19) can be transformed into the unconstrained problem:

$$\min_{\mathbf{z} \in \mathbb{R}} (1/2) \|\mathbf{Uz} + \bar{\mathbf{x}} - \mathbf{s}\|_2^2.$$

This problem has a closed form solution $\mathbf{z}^* = (\mathbf{U}^T \mathbf{U})^{-1} \mathbf{U}^T (\mathbf{s} - \bar{\mathbf{x}}) = \mathbf{U}^T (\mathbf{s} - \bar{\mathbf{x}})$.

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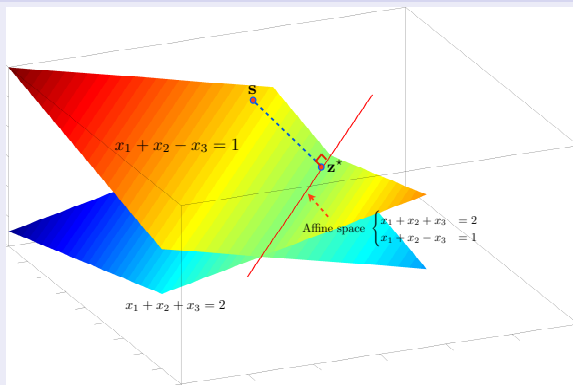
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Illustration



Limitations of the null-space method

Limitations of the null space approach

- ▶ Require **matrix factorization** (e.g., QR factorization) to compute a basis \mathbf{U} of the **null space** of \mathbf{A} and a **feasible point** $\bar{\mathbf{x}}$, which is computational demand in high-dimension ($\mathcal{O}(n^2p)$).
- ▶ If matrix \mathbf{A} is given **implicitly** (e.g., by linear operator), then computing \mathbf{U} is **impractical**.
- ▶ Null space method **destroys** the original structure of the objective function f due to the **affine transformation** $\mathbf{U}\mathbf{z} + \bar{\mathbf{x}}$. For instance, $f(\mathbf{x}) := \|\mathbf{x}\|_1$, which is component-wise decomposable.

Convex problems with simple constraints

Convex problems with simple constraints

When $\mathbf{Ax} = \mathbf{b}$ is **absent**, problem (1) reduces to:

$$f^* := \min_{\mathbf{x} \in \mathcal{X}} f(\mathbf{x}) \quad (20)$$

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Assumption (Simplicity)

\mathcal{X} is “**simple**” so that the **projection** $\pi_{\mathcal{X}}$ of any point $\mathbf{s} \in \mathbb{R}^p$ onto \mathcal{X} can be computed **efficiently**, i.e.:

$$\pi_{\mathcal{X}}(\mathbf{s}) := \arg \min_{\mathbf{x} \in \mathcal{X}} \|\mathbf{x} - \mathbf{s}\|_2,$$

can be solved **efficiently** (e.g., **closed form solution** or **polynomial time**).

Note: Let $\iota_{\mathcal{X}}$ be the **indicator function** of \mathcal{X} . Then

$$\pi_{\mathcal{X}}(\mathbf{s}) = \text{prox}_{\iota_{\mathcal{X}}}(\mathbf{s}).$$

Examples can be found in Lectures 4 and 5.

Projected-gradient method

Assumption A.1

- ▶ $f \in \mathcal{F}_L^{1,1}(\mathbb{R}^p)$
- ▶ $\pi_{\mathcal{X}}$ can be computed exactly.

Projected-gradient method

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Projected gradient method (ProjGA)

1. Choose $\mathbf{x}^0 \in \mathbb{R}^p$.
2. For $k = 0, 1, \dots$, perform:

$$\mathbf{x}^{k+1} := \pi_{\mathcal{X}}(\mathbf{x}^k - (1/L_f)\nabla f(\mathbf{x}^k)).$$

Projected-gradient method

Assumption A.1

- ▶ $f \in \mathcal{F}_L^{1,1}(\mathbb{R}^p)$
- ▶ $\pi_{\mathcal{X}}$ can be **computed exactly**.

Projected gradient method (ProjGA)

1. Choose $\mathbf{x}^0 \in \mathbb{R}^p$.
2. For $k = 0, 1, \dots$, perform:

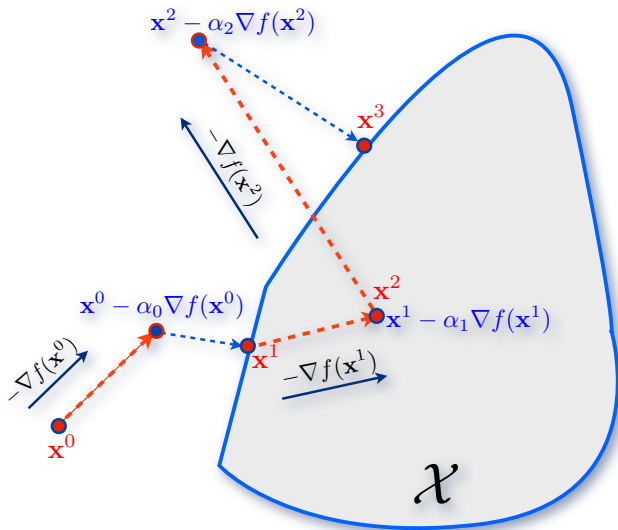
$$\mathbf{x}^{k+1} := \pi_{\mathcal{X}}(\mathbf{x}^k - (1/L_f)\nabla f(\mathbf{x}^k)).$$

Properties

- ▶ ProjGA can be enhanced by performing a **line-search** for approximating L_f .
- ▶ **Convergence**: The convergence of ProjGA remains the same as in **standard gradient method**, i.e.:

$$f(\mathbf{x}^k) - f^* \leq \frac{L_f \|\mathbf{x}^0 - \mathbf{x}^*\|_2^2}{2(k+1)}, \quad k \geq 0.$$

Illustration of the projected gradient method



Three iterations of the projected gradient method.

Fast projected-gradient method

Assumption

Under **Assumption A.1.**, ProjGA can be **accelerated** by using **Nesterov's optimal method**.

Fast projected gradient method (FastProjGA)

1. Choose $\mathbf{x}^0 \in \mathbb{R}^p$. Set $\mathbf{y}^0 := \mathbf{x}^0$ and $t_0 := 1$
2. For $k = 0, 1, \dots$, perform:

$$\begin{cases} \mathbf{x}^{k+1} & := \pi_{\mathcal{X}}(\mathbf{y}^k - (1/L_f)\nabla f(\mathbf{y}^k)), \\ \mathbf{y}^{k+1} & := \mathbf{x}^{k+1} + ((t_k - 1)/t_{k+1})(\mathbf{x}^{k+1} - \mathbf{x}^k), \\ t_{k+1} & := (1 + \sqrt{1 + 4t_k^2})/2. \end{cases}$$

Fast projected-gradient method

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Under **Assumption A.1.**, ProjGA can be **accelerated** by using **Nesterov's optimal method**.

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Convergence

The convergence of FastProjGA remains the same as in **fast gradient method**, i.e.:

$$f(\mathbf{x}^k) - f^* \leq \frac{2L_f \|\mathbf{x}^0 - \mathbf{x}^*\|_2^2}{(k+1)^2}, \quad k \geq 0.$$

Frank-Wolfe's method

Problem setting and assumption

$$f^* := \min_{\mathbf{x} \in \mathcal{X}} f(\mathbf{x}) \quad (21)$$

Assumptions

- ▶ \mathcal{X} is nonempty, **convex**, closed and **bounded**.
- ▶ $f \in \mathcal{F}_L^{1,1}(\mathbb{R}^p)$ (i.e., convex with Lipschitz gradient).
- ▶ For given $c \in \mathbb{R}^p$, $\hat{\mathbf{x}} := \arg \min_{\mathbf{x} \in \mathcal{X}} c^T \mathbf{x}$ can be solved **efficiently**.

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Frank-Wolfe's method [5]

Conditional gradient method (CGA)

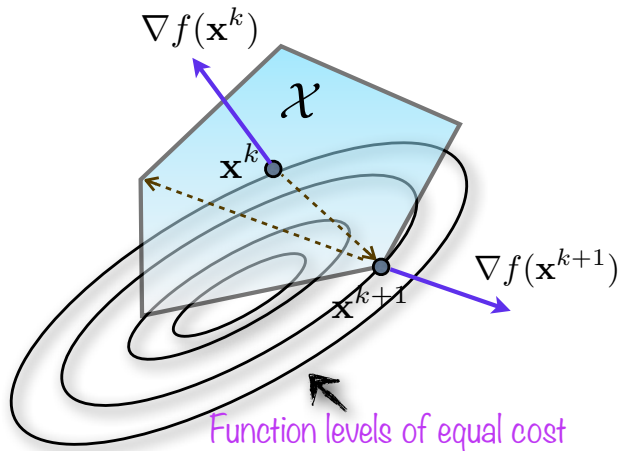
1. Choose $\mathbf{x}^0 \in \mathcal{X}$.
2. For $k = 0, 1, \dots$, perform:

$$\begin{cases} \hat{\mathbf{x}}^k & := \arg \min_{\mathbf{x} \in \mathcal{X}} \nabla f(\mathbf{x}^k)^T \mathbf{x}, \\ \mathbf{x}^{k+1} & := (1 - \gamma_k) \mathbf{x}^k + \gamma_k \hat{\mathbf{x}}^k, \end{cases}$$

where $\gamma_k := \frac{2}{k+2}$ is a given relaxation parameter.

Geometric interpretation of Frank-Wolfe's method

- ▶ Most straightforward way to generate a feasible *descent direction*: find $\hat{\mathbf{x}}^k$ that satisfies $\nabla f(\mathbf{x}^k)^T(\hat{\mathbf{x}}^k - \mathbf{x}^k) < 0$.
- ▶ We assume that the constraint set \mathcal{X} is compact so that the direction finding problem has a solution.



Properties and convergence of Frank-Wolfe's method

Properties

- ▶ Since \mathcal{X} is **bounded**, \hat{x}^k is **well-defined**.
- ▶ CGA is a “**norm-free**” method
- ▶ \hat{x}^k attains at the **boundary** of \mathcal{X} , which **preserves sparsity**.
- ▶ When \mathcal{X} is a **polytope**, computing \hat{x}^k is equivalent to solving a **linear program**.
- ▶ Allows inexactness in computing \hat{x}^k
- ▶ γ_k can be estimated by a line-search procedure.

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- ▶ Allows inexactness in computing $\hat{\mathbf{x}}^k$
- ▶ γ_k can be estimated by a line-search procedure.

Theorem (Convergence [5])

Let $\{\mathbf{x}^k\}$ be the sequence generated by CGA. Then

$$f(\mathbf{x}^k) - f^* \leq \frac{2L_f}{k+1} D_{\mathcal{X}}^2,$$

where $D_{\mathcal{X}} := \max_{\mathbf{x}, \mathbf{y} \in \mathcal{X}} \|\mathbf{x} - \mathbf{y}\|$, the diameter of \mathcal{X} w.r.t. $\|\cdot\|$.

The **convergence rate** of **CGA** is $\mathcal{O}(1/k)$ which is the **same order** as **ProjGA**. However, the **diameter** $D_{\mathcal{X}}$ is in general **worse** than $\|\mathbf{x}^0 - \mathbf{x}^*\|_2$ in **ProjGA** in the ℓ_2 -norm.

Dual subgradient method

Dual problem (13) is in general **nonsmooth and convex**. **Subgradient ascent method** can be applied to solve it.

Properties of dual function

- ▶ d is **concave**, but **not necessary differentiable**.
- ▶ **Subgradient:** $\mathbf{Ax}^*(\lambda) - \mathbf{b} \in \partial d(\lambda)$, where $\mathbf{x}^*(\lambda)$ is a solution of (12).

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Dual subgradient ascent method

Dual subgradient method (DSGM):

1. Choose $\lambda^0 \in \mathbb{R}^p$.
2. For $k = 0, 1, \dots$, perform:
 - 2.a. Solve (12) to obtain $\mathbf{x}^*(\lambda)$.
 - 2.b. Compute the **subgradient** $\nabla d(\lambda^k) := \mathbf{Ax}^*(\lambda^k) - \mathbf{b}$.
 - 2.c. Update $\lambda^{k+1} := \lambda^k + \frac{R}{\sqrt{k+1}} \nabla d(\lambda^k)$, where R is a given constant.

Convergence of DSGM

Well-definedness

- ▶ Problem (12) may not have solution $\mathbf{x}^*(\lambda)$ for any λ . Then DSGM is not well-defined except \mathcal{X} is bounded.
- ▶ Impractical to evaluate $R_\star := \|\lambda^0 - \lambda^\star\|_2$, use an upper bound R of R_\star .

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- ▶ **Impractical** to evaluate $R_\star := \|\lambda^0 - \lambda^\star\|_2$, use an **upper bound** R of R_\star .

Theorem (Convergence)

Assume that $\|\mathbf{Ax}^\star(\lambda^k) - \mathbf{b}\| \leq M_d$ for all $k \geq 0$. Then $\{\lambda^k\}$ generated by DSGM satisfies

$$d^\star - d(\lambda^k) \leq \frac{M_d R_\star}{\sqrt{k+1}}, \forall k \geq 0,$$

where $R_\star := \min_{\lambda^\star} \|\lambda^0 - \lambda^\star\|_2$. *Convergence rate of DSGM is $\mathcal{O}(1/\sqrt{k})$.*

Convergence of DSGM

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Special cases

1. If both f is **strongly convex**, then d is **smooth** and its **gradient** is **Lipschitz continuous**., $d \in \mathcal{F}_L^{1,1}(\mathbb{R}^p)$. **Gradient** and **fast gradient methods** in Lecture 3 can be used to solve the **dual problem**.
2. **Smoothing techniques** in Lecture 5 can be used to smooth the dual function d .

Augmented Lagrangian method

Dual problem (13) is **convex** but generally **nonsmooth**. By augmenting \mathcal{L} with $(\kappa/2)\|\mathbf{Ax} - \mathbf{b}\|_2^2$, we obtain **augmented dual function** d_κ , which maintains basic properties of d but **smooth and Lipschitz gradient**.

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Augmented Lagrangian and augmented dual function

- ▶ **Augmented Lagrangian:** $\mathcal{L}_\kappa(\mathbf{x}, \lambda) := \mathcal{L}(\mathbf{x}, \lambda) + (\kappa/2)\|\mathbf{Ax} - \mathbf{b}\|_2^2$, where $\rho > 0$ is a penalty parameter.
- ▶ **Augmented dual function:**

$$d_\kappa(\lambda) := \min_{\mathbf{x} \in \mathcal{X}} \left\{ \mathcal{L}_\kappa(\mathbf{x}, \lambda) := f(\mathbf{x}) + \lambda^T(\mathbf{Ax} - \mathbf{b}) + (\kappa/2)\|\mathbf{Ax} - \mathbf{b}\|_2^2 \right\}. \quad (22)$$

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Key properties of d_κ

- ▶ d_κ is **concave** and **smooth** and

$$\nabla d_\kappa(\lambda) = \mathbf{Ax}_\kappa^*(\lambda) - \mathbf{b},$$

where $\mathbf{x}_\kappa^*(\lambda)$ is the **solution** of (22).

- ▶ ∇d_κ is **Lipschitz continuous** with a Lipschitz constant $L_d := \kappa^{-1}$, i.e.:

$$\|\nabla d_\kappa(\lambda) - \nabla d_\kappa(\hat{\lambda})\| \leq \kappa^{-1}\|\lambda - \hat{\lambda}\|, \quad \forall \lambda, \hat{\lambda} \in \mathbb{R}^n.$$

Example: Behavior of the augmented Lagrangian dual function

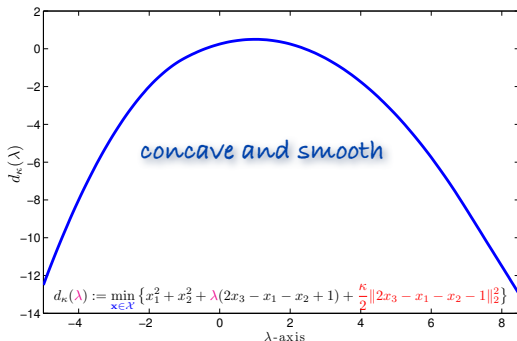
Consider a constrained convex problem:

$$\begin{aligned} \min_{\mathbf{x} \in \mathbb{R}^3} \quad & \{f(\mathbf{x}) := x_1^2 + x_2^2\}, \\ \text{s.t.} \quad & 2x_3 - x_1 - x_2 = 1, \\ & \mathbf{x} \in \mathcal{X} := [-2, 2] \times [-2, 2] \times [0, 2]. \end{aligned}$$

The augmented Lagrangian dual function is defined as

$$d_\kappa(\lambda) := \min_{\mathbf{x} \in \mathcal{X}} \{x_1^2 + x_2^2 + \lambda(2x_3 - x_1 - x_2 + 1) + (\kappa/2)\|2x_3 - x_1 - x_2 - 1\|_2^2\}$$

is **concave** and **nonsmooth** as illustrated in the figure below.



Augmented dual problem

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$$d_{\kappa}^* := \max_{\lambda \in \mathbb{R}^n} d_{\kappa}(\lambda), \quad \kappa > 0. \quad (23)$$

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Relation to the dual problem (13)

Under **Slater's condition** and $\mathcal{X}^* \neq \emptyset$, we have

- ▶ The **dual solution set** of (23) is coincided with the **one** of the **dual problem (13)**.
- ▶ $f^* = d^* = d_{\kappa}^*$ for any $\kappa > 0$.

The **augmented dual problem (23)** is **smooth and convex** \Rightarrow **Gradient and Fast gradient methods** can be applied to solve it.

Augmented Lagrangian method

Augmented Lagrangian method (ALM):

1. Choose $\lambda^0 \in \mathbb{R}^p$ and $\kappa > 0$.
2. For $k = 0, 1, \dots$, perform:
 - 2.a. Solve (22) to compute $\nabla d_\kappa(\lambda^k) := \mathbf{A}x_\kappa^*(\lambda^k) - \mathbf{b}$.
 - 2.b. Update $\lambda^{k+1} := \lambda^k + \kappa \nabla d_\kappa(\lambda^k)$.

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ALM can be accelerated by **Nesterov's optimal method**.

Fast augmented Lagrangian method (FALM)

1. Choose $\lambda^0 \in \mathbb{R}^p$ and $\kappa > 0$. Set $\tilde{\lambda}^0 := \lambda^0$ and $t_0 := 1$.
2. For $k = 0, 1, \dots$, perform:
 - 2.a. Solve (22) to compute $\nabla d_\kappa(\tilde{\lambda}^k) := \mathbf{A}\mathbf{x}_\kappa^*(\tilde{\lambda}^k) - \mathbf{b}$.
 - 2.b. Update

$$\begin{cases} \lambda^{k+1} & := \tilde{\lambda}^k + \kappa \nabla d_\kappa(\tilde{\lambda}^k), \\ \tilde{\lambda}^{k+1} & := \lambda^{k+1} + ((t_k - 1)/t_{k+1})(\lambda^{k+1} - \lambda^k), \\ t_{k+1} & := (1 + \sqrt{1 + 4t_k^2})/2. \end{cases}$$

Convergence of ALM and FALM

Theorem (Convergence)

- ▶ Let $\{\lambda^k\}$ be the sequence generated by ALM. Then

$$d^* - d_\kappa(\lambda^k) \leq \frac{\|\lambda^0 - \lambda^*\|_2^2}{2\kappa(k+1)}, \quad k \geq 0.$$

- ▶ Let $\{\lambda^k\}$ be the sequence generated by FALM. Then

$$d^* - d_\kappa(\lambda^k) \leq \frac{2\|\lambda^0 - \lambda^*\|_2^2}{\kappa(k+2)^2}, \quad k \geq 0.$$

- ▶ The **convergence rate** of ALM is $\mathcal{O}(1/k)$ w.r.t. the augmented dual function d_κ .
- ▶ The **convergence rate** of FALM is $\mathcal{O}(1/k^2)$ w.r.t. the augmented dual function d_κ .
- ▶ **Important observation:** The right-hand side of both estimates **depends on κ** . When κ is getting **large**, the right-hand side is **decreasing**.

Drawbacks and enhancements

Drawbacks

1. **Drawback 1:** The quadratic term $\|\mathbf{Ax} - \mathbf{b}\|_2^2$ in (22) **destroys** the **separability** as well as the **tractable proximity** of f .
2. **Drawback 2:** Solving (22) exactly is **impractical**.
3. **Drawback 3:** **No theoretical guarantee** for choosing appropriate values of κ .

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3. **Drawback 3:** **No theoretical guarantee** for choosing appropriate values of κ .

Enhancements

1. Allow **inexactness** of solving (22), while guaranteeing the **same convergence rate**.
2. Update the penalty parameter κ
 - ▶ **Increasing ρ :** Lead to the increase of ill-condition in (22).
 - ▶ **Adaptively update κ :** Often heuristic
3. Process the quadratic term $\|\mathbf{Ax} - \mathbf{b}\|_2^2$ by linearization, alternating, etc.

Example: Group basis pursuit

Group basis pursuit

Given a linear operator \mathbf{A} , a measurement vector \mathbf{b} and a group structure $\mathcal{G} := \{\mathcal{G}_1, \dots, \mathcal{G}_g\}$. The aim is to solve:

$$\min_{\mathbf{x} \in \mathbb{R}^p} \sum_{i=1}^g \|\mathbf{x}_{\mathcal{G}_i}\|_2 \quad \text{s.t.} \quad \mathbf{Ax} = \mathbf{b}. \quad (24)$$

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Applying ALM and FALM

The main computation:

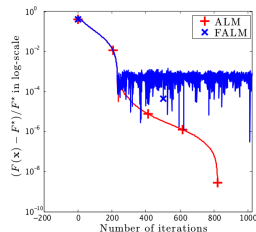
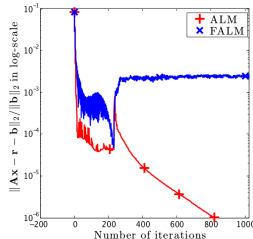
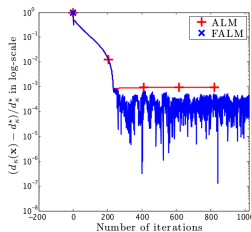
- ▶ Solving the subproblem (22), which is

$$\mathbf{x}_\kappa^*(\lambda) := \arg \min_{\mathbf{x} \in \mathcal{X}} \left\{ \sum_{i=1}^g \|\mathbf{x}_{\mathcal{G}_i}\|_2 + \lambda^T (\mathbf{Ax} - \mathbf{b}) + (\kappa/2) \|\mathbf{Ax} - \mathbf{b}\|_2^2 \right\},$$

by applying, e.g., FISTA (Lecture 5).

- ▶ Updating κ by increasing it as $\kappa_{k+1} := \eta \kappa_k$ for given $\eta > 1$.

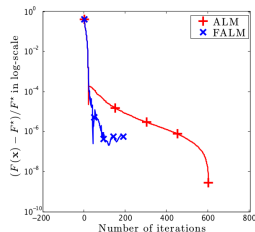
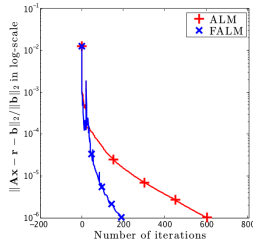
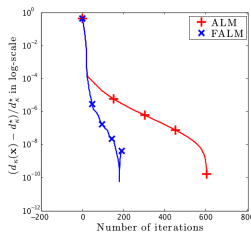
Numerical results



	ALM	FALM
Primal Obj. Value	47.145	47.187
Feas. Gap	0.99×10^{-6}	0.23×10^{-2}
Dual Obj. Value	33.196	33.165
Iterations	821	2000
CPU time (s)	2.656	6.513
Calls A/A^T	9031/8210	22000/20000
Recovery error	0.04%	0.4%

- Parameters: $\kappa = 0.5$, $\eta = 1$
- Input: $n = 341$, $p = 1024$, $g = 85$, $\text{nzg} = 11$; $\min |\mathcal{G}_i| = 5$, $\max |\mathcal{G}_i| = 23$, $\text{mean} |\mathcal{G}_i| = 12.04$
- Proximal operations (FISTA): max iterations 10, stop criteria 10^{-9} relative change, warm start
- Stopping criteria: $\|\mathbf{Ax}^k - \mathbf{r}^k - \mathbf{b}\| \leq 10^{-6} \|\mathbf{b}\|$ and $\|(\mathbf{x}^k, \mathbf{r}^k) - (\mathbf{x}^{k-1}, \mathbf{r}^{k-1})\| \leq 10^{-6} \|(\mathbf{x}^k, \mathbf{r}^k)\|$

Numerical results



	ALM	FALM
Primal Obj. Value	47.1451	47.1452
Feas. Gap	0.99×10^{-6}	0.99×10^{-6}
Dual Obj. Value	33.196	33.196
Iterations	605	192
CPU time (s)	10.647	4.920
Calls A/A^T	38348/37743	17420/17228
Recovery error	0.04%	0.04%

- Parameters: $\kappa = 0.5$, $\eta = 1$
- Input: $n = 341$, $p = 1024$, $g = 85$, $\text{nzg} = 11$; $\min |\mathcal{G}_i| = 5$, $\max |\mathcal{G}_i| = 23$, $\text{mean} |\mathcal{G}_i| = 12.04$
- Proximal operations (FISTA): max iterations 100, stop criteria 10^{-9} relative change, warm start
- Stopping criteria: $\|\mathbf{Ax}^k - \mathbf{r}^k - \mathbf{b}\| \leq 10^{-6} \|\mathbf{b}\|$ and $\|(\mathbf{x}^k, \mathbf{r}^k) - (\mathbf{x}^{k-1}, \mathbf{r}^{k-1})\| \leq 10^{-6} \|(\mathbf{x}^k, \mathbf{r}^k)\|$

Remarks

Remarks

- ▶ The FALM method is sensitive to the inexactness of the solution of (22)

$$\mathbf{x}_\kappa^*(\lambda) := \arg \min_{\mathbf{x} \in \mathcal{X}} \left\{ \sum_{i=1}^g \|\mathbf{x}_{\mathcal{G}_i}\|_2 + \lambda^T (\mathbf{A}\mathbf{x} - \mathbf{b}) + (\kappa/2) \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2 \right\}$$

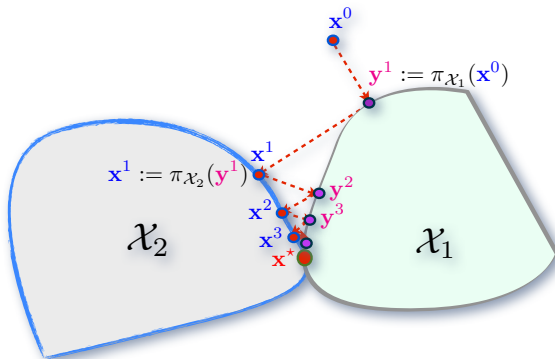
- ▶ "Fast" updates of the dual variable λ^k influence the primal updates
 - ▶ warm-start strategy - at iteration k choose initial solution of (22) $x_\kappa^*(\lambda^{k-1})$
 - ▶ increase iterations number to achieve convergence of the primal (also tolerance)
 - ▶ keep η small (FALM more sensitive to large values of η)
- ▶ Guarantes are given only for the dual problem, not for the primal

Alternating idea to overcome the non-separability

- ▶ **Problem:** Given two nonempty, closed and convex sets \mathcal{X}_1 and \mathcal{X}_2 . Find a point $\mathbf{x}^* \in \mathcal{X}_1 \cap \mathcal{X}_2$.
- ▶ **Strategy:** Start from \mathbf{x}^0 and iterate alternatively:

$$\begin{cases} \mathbf{y}^{k+1} & := \pi_{\mathcal{X}_1}(\mathbf{x}^k) \\ \mathbf{x}^{k+1} & := \pi_{\mathcal{X}_2}(\mathbf{y}^{k+1}) \end{cases}$$

where $\pi_{\mathcal{X}}$ is the projection on the convex set \mathcal{X} .



Alternating minimization algorithm (AMA)

Assumptions

- ▶ Problem (1) has a **separable structure** with $p = 2$, i.e.:

$$f^* := \begin{cases} \min_{\mathbf{x} \in \mathbb{R}^p} & \{f(\mathbf{x}) := f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2)\}, \\ \text{s.t.} & \mathbf{A}_1 \mathbf{x}_1 + \mathbf{A}_2 \mathbf{x}_2 = \mathbf{b}, \mathbf{x}_1 \in \mathcal{X}_1, \mathbf{x}_2 \in \mathcal{X}_2. \end{cases} \quad (25)$$

- ▶ f_1 is **strongly convex** with parameter $\mu_1 > 0$.

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Assumptions

- ▶ Problem (1) has a **separable structure** with $p = 2$, i.e.:

$$f^* := \begin{cases} \min_{\mathbf{x} \in \mathbb{R}^p} & \{f(\mathbf{x}) := f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2)\}, \\ \text{s.t.} & \mathbf{A}_1 \mathbf{x}_1 + \mathbf{A}_2 \mathbf{x}_2 = \mathbf{b}, \mathbf{x}_1 \in \mathcal{X}_1, \mathbf{x}_2 \in \mathcal{X}_2. \end{cases} \quad (25)$$

- ▶ f_1 is **strongly convex** with parameter $\mu_1 > 0$.

The idea of AMA [7]

- ▶ **Alternating** between variables \mathbf{x}_1 and \mathbf{x}_2 in:

$$\min_{\mathbf{x}_1 \in \mathcal{X}_1, \mathbf{x}_2 \in \mathcal{X}_2} \left\{ f_1(\mathbf{x}_1) + f_2(\mathbf{x}_2) + \lambda^T \mathbf{A}_1 \mathbf{x}_1 + \lambda^T \mathbf{A}_2 \mathbf{x}_2 + (\kappa/2) \|\mathbf{A}_1 \mathbf{x}_1 + \mathbf{A}_2 \mathbf{x}_2 - \mathbf{b}\|_2^2 \right\}.$$

- ▶ Since f_1 is **convex**, **neglects** the augmented term. Then, this step becomes

$$\begin{cases} \mathbf{x}_1^{k+1} & := \arg \min_{\mathbf{x}_1 \in \mathcal{X}_1} \left\{ f_1(\mathbf{x}_1) + (\lambda^k)^T \mathbf{A}_1 \mathbf{x}_1 \right\}, \\ \mathbf{x}_2^{k+1} & := \arg \min_{\mathbf{x}_2 \in \mathcal{X}_2} \left\{ f_2(\mathbf{x}_2) + (\lambda^k)^T \mathbf{A}_2 \mathbf{x}_2 + \frac{\kappa}{2} \|\mathbf{A}_1 \mathbf{x}_1^{k+1} + \mathbf{A}_2 \mathbf{x}_2 - \mathbf{b}\|_2^2 \right\}. \end{cases}$$

AMA: Alternating minimization algorithm

Alternating minimization algorithm (AMA):

1. Choose $\lambda^0 \in \mathbb{R}^p$ and $\kappa > 0$.

2. For $k = 0, 1, \dots$, perform:

$$\begin{cases} \mathbf{x}_1^{k+1} & := \arg \min_{\mathbf{x}_1 \in \mathcal{X}_1} \{f_1(\mathbf{x}_1) + (\lambda^k)^T \mathbf{A}_1 \mathbf{x}_1\} \\ \mathbf{x}_2^{k+1} & := \arg \min_{\mathbf{x}_2 \in \mathcal{X}_2} \left\{ f_2(\mathbf{x}_2) + (\lambda^k)^T \mathbf{A}_2 \mathbf{x}_2 + \frac{\kappa}{2} \|\mathbf{A}_1 \mathbf{x}_1^{k+1} + \mathbf{A}_2 \mathbf{x}_2 - \mathbf{b}\|_2^2 \right\} \\ \lambda^{k+1} & := \lambda^k + \kappa (\mathbf{A}_1 \mathbf{x}_1^{k+1} + \mathbf{A}_2 \mathbf{x}_2^{k+1} - \mathbf{b}). \end{cases}$$

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Implementation remarks

- ▶ **Main computation:** Solving two subproblems to compute \mathbf{x}_1^{k+1} and \mathbf{x}_2^{k+1} .
- ▶ \mathbf{A}_2 prevents the tractable proximity from f_2 .
- ▶ When $\mathbf{A}_2^T \mathbf{A}_2 = \mathbf{I}$, we have $\mathbf{x}_2^{k+1} = \text{prox}_{\kappa^{-1} f_2}(\mathbf{A}_2^T (\mathbf{b} - \mathbf{A}_1 \mathbf{x}_1^{k+1}) - \kappa^{-1} \mathbf{A}_2^T \lambda^k)$.
- ▶ When $\mathbf{A}_2^T \mathbf{A}_2 \neq \mathbf{I}$, we can approximate \mathbf{x}_2^{k+1} by linearizing the quadratic term.
- ▶ The penalty parameter κ can be updated.

Convergence of AMA

Observations

- ▶ AMA is a **proximal-gradient method** applying to the Fenchel **dual problem**:

$$\tilde{d}^* := \max_{\lambda \in \mathbb{R}^p} \left\{ \tilde{d}(\lambda) := -f_1^*(-\mathbf{A}_1^T \lambda) - f_2^*(-\mathbf{A}_2^T \lambda) - \mathbf{b}^T \lambda \right\}. \quad (26)$$

where f_1^* and f_2^* are the **Fenchel conjugate** of f_1 and f_2 , respectively.

- ▶ Since f_1 is **strongly convex**, the conjugate f_1^* is **Lipschitz gradient** with Lipschitz constant $L_{f_1^*} := \mu_1^{-1}$.
- ▶ AMA can be **accelerated** by using **Nesterov's optimal gradient method** (see [3]).

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- AMA can be **accelerated** by using **Nesterov's optimal gradient method** (see [3]).

Theorem (Convergence theorem [3])

Let $\{(\mathbf{x}_1^k, \mathbf{x}_2^k, \lambda^k)\}$ be the sequence generated by AMA. Assume that $\rho < 2\mu_1 / \lambda_{\max}(\mathbf{A}_1^T \mathbf{A}_1)$. Then

$$\tilde{d}^* - \tilde{d}(\lambda^k) \leq \frac{\lambda_{\max}(\mathbf{A}_1^T \mathbf{A}_1)}{2\mu_1(k+1)} \|\lambda^0 - \lambda^*\|_2^2,$$

where $\lambda_{\max}(\mathbf{A}_1^T \mathbf{A}_1)$ is the maximum eigenvalue of $\mathbf{A}_1^T \mathbf{A}_1$.

Example: ℓ_1 -regularized least squaresProblem (ℓ_1 -regularized least squares)

$$\min_{\mathbf{x} \in \mathbb{R}^p} (1/2) \|\mathbf{Ax} - \mathbf{b}\|_2^2 + \rho \|\mathbf{x}\|_1, \quad (27)$$

where $\rho > 0$ is a *regularization parameter*.

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Problem (ℓ_1 -regularized least squares)

$$\min_{\mathbf{x} \in \mathbb{R}^p} (1/2) \|\mathbf{Ax} - \mathbf{b}\|_2^2 + \rho \|\mathbf{x}\|_1, \quad (27)$$

where $\rho > 0$ is a *regularization parameter*.

Applying AMA

Introducing a **slack variable** $\mathbf{r} = \mathbf{Ax} - \mathbf{b}$, we can reformulate (27) as

$$\min_{\mathbf{x} \in \mathbb{R}^p, \mathbf{r} \in \mathbb{R}^n} (1/2) \|\mathbf{r}\|_2^2 + \rho \|\mathbf{x}\|_1, \quad \text{s.t. } \mathbf{Ax} - \mathbf{r} = \mathbf{b}.$$

The main steps of AMA becomes

$$\begin{cases} \mathbf{r}^{k+1} & := \arg \min_{\mathbf{r} \in \mathbb{R}^n} \left\{ (1/2) \|\mathbf{r}\|_2^2 - (\lambda^k)^T \mathbf{r} \right\} \equiv \lambda^k \\ \mathbf{x}^{k+1} & := \arg \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ \rho \|\mathbf{x}\|_1 + (\lambda^k)^T \mathbf{Ax} + \frac{\kappa}{2} \|\mathbf{Ax} - \mathbf{r}^{k+1} - \mathbf{b}\|_2^2 \right\}, \\ \lambda^{k+1} & := \lambda^k + \kappa (\mathbf{Ax}^{k+1} - \mathbf{r}^{k+1} - \mathbf{b}). \end{cases}$$

For $\mathbf{A}^T \mathbf{A} = \mathbb{I}$, the \mathbf{x} -step reduces to:

$$\mathbf{x}^{k+1} := \text{prox}_{\kappa^{-1} \rho \|\cdot\|_1} \left(\mathbf{A}^T (\mathbf{b} + \lambda^k) - \kappa^{-1} \mathbf{A}^T \lambda^k \right).$$

Approaches to solving the subproblem

Problem

- ▶ The main computation of AMA is the solution of:

$$\mathbf{x}^{k+1} := \arg \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ \rho \|\mathbf{x}\|_1 + (\lambda^k)^T \mathbf{A} \mathbf{x} + \frac{\kappa}{2} \|\mathbf{A} \mathbf{x} - \mathbf{r}^{k+1} - \mathbf{b}\|_2^2 \right\} \quad (28)$$

- ▶ (28) has no closed form solution (except for $\mathbf{A}^T \mathbf{A} = \mathbb{I}$).

Solution

- ▶ There are two ways to overcome this drawback:
 - ▶ Applying FISTA.
 - ▶ Linearize the quadratic term: $q(\mathbf{x}) := q(\mathbf{x}^k) + \nabla q(\mathbf{x}^k)^T (\mathbf{x} - \mathbf{x}^k) + \frac{L}{2} \|\mathbf{x} - \mathbf{x}^k\|_2^2$ where L is the Lipschitz constant equal to $\|\mathbf{A}\|_2^2$

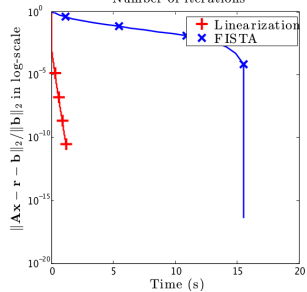
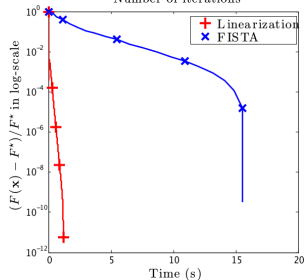
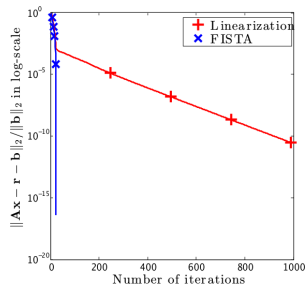
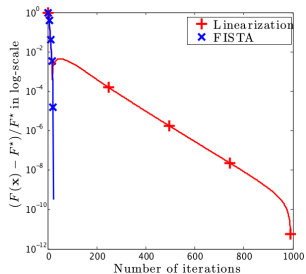
Note: Is equivalent to applying FISTA with 1 iteration

Numerical results - High accuracy

	Linearization	FISTA
Primal Obj. Value	14.241	14.241
Feas. Gap	0.3×10^{-10}	0.3×10^{-17}
Iterations	991	23
Inner Iterations	991	13835
CPU time (s)	1.187	15.555
Calls A/A^T	992/991	13859/13835

- Parameters: $\rho = 0.1$, $\kappa = 0.01$, $\eta = 1.25$
- Input: $n = 750$, $p = 2000$, $k = 200$, Noise $\sim \mathcal{N}(0, \sigma^2 \mathcal{I})$ with $\sigma = 10^{-3}$
- FISTA: max iterations **1000**, stop criteria 10^{-10} relative change, warm start
- Stopping criteria: $\|\mathbf{Ax}^k - \mathbf{r}^k - \mathbf{b}\| \leq 10^{-10} \|\mathbf{b}\|$ and $\|\mathbf{x}^k - \mathbf{x}^{k-1}\| \leq 10^{-10} \|\mathbf{x}^k\|$

Convergence plots

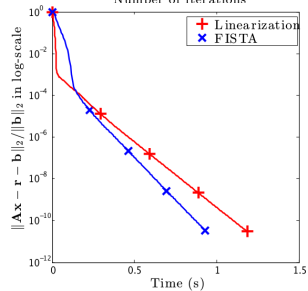
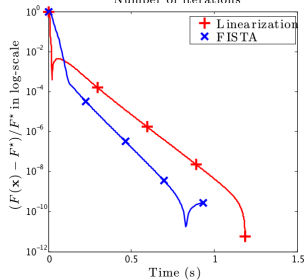
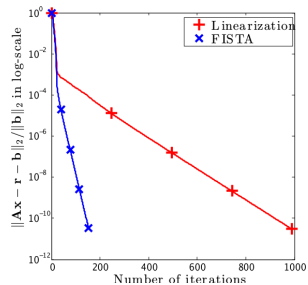
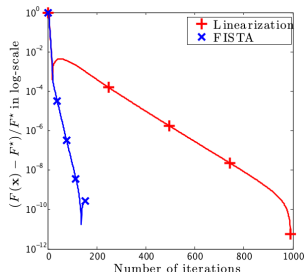


Numerical results - Low accuracy

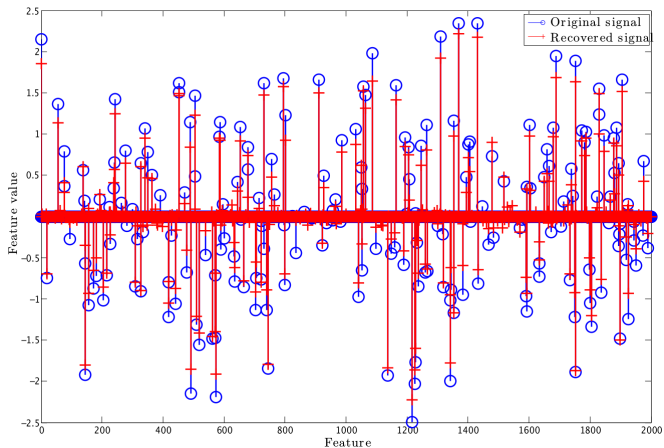
	Linearization	FISTA
Primal Obj. Value	14.241	14.241
Feas. Gap	0.3×10^{-10}	0.29×10^{-10}
Iterations	991	154
Inner Iterations	991	758
CPU time (s)	1.187	0.938
Calls A/A^T	992/991	913/758

- ▶ Parameters: $\rho = 0.1$, $\kappa = 0.01$, $\eta = 1.25$
- ▶ Input: $n = 750$, $p = 2000$, $k = 200$, Noise $\sim \mathcal{N}(0, \sigma^2 \mathcal{I})$ with $\sigma = 10^{-3}$
- ▶ FISTA: max iterations **5**, stop criteria 10^{-10} relative change, warm start
- ▶ Stopping criteria: $\|\mathbf{Ax}^k - \mathbf{r}^k - \mathbf{b}\| \leq 10^{-10} \|\mathbf{b}\|$ and $\|\mathbf{x}^k - \mathbf{x}^{k-1}\| \leq 10^{-10} \|\mathbf{x}^k\|$

Convergence plots



Recovery error



- ▶ $\|\mathbf{x}^* - \mathbf{x}^{\natural}\| / \|\mathbf{x}^{\natural}\|$
 - ▶ Linearization: 18.88%
 - ▶ FISTA: 18.88%
- ▶ $\|\mathbf{x}_{\text{Lin}}^* - \mathbf{x}_{\text{FISTA}}^*\| / \|\mathbf{x}^{\natural}\| = 0.43 \times 10^{-8}$

Alternating direction method of multipliers (ADMM)

The idea

When f_1 is **not strongly convex**, to overcome the drawback of ALM, by **alternating** solving (22).

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ADMM

Alternating direction method of multipliers (ADMM):

1. Choose $\lambda^0 \in \mathbb{R}^p$, $\mathbf{x}_1^0 \in \mathbb{R}^p$, $\gamma \geq 0$ and $\kappa > 0$.
2. For $k = 0, 1, \dots$, perform:

$$\begin{cases} \mathbf{x}_1^{k+1} := \operatorname{argmin}_{\mathbf{x}_1 \in \mathcal{X}_1} \left\{ f_1(\mathbf{x}_1) + \frac{\kappa}{2} \|\mathbf{A}_1 \mathbf{x}_1 + \mathbf{A}_2 \mathbf{x}_2^k - \mathbf{b} - \kappa^{-1} \mathbf{A}_1^T \lambda^k\|_2^2 + \frac{\gamma}{2} \|\mathbf{x}_1 - \mathbf{x}_1^k\|_2^2 \right\}, \\ \mathbf{x}_2^{k+1} := \operatorname{argmin}_{\mathbf{x}_2 \in \mathcal{X}_2} \left\{ f_2(\mathbf{x}_2) + \frac{\kappa}{2} \|\mathbf{A}_1 \mathbf{x}_1^{k+1} + \mathbf{A}_2 \mathbf{x}_2 - \mathbf{b} - \kappa^{-1} \mathbf{A}_2^T \lambda^k\|_2^2 \right\}, \\ \lambda^{k+1} := \lambda^k + \kappa (\mathbf{A}_1 \mathbf{x}_1^{k+1} + \mathbf{A}_2 \mathbf{x}_2^{k+1} - \mathbf{b}). \end{cases}$$

In the original ADMM version, the **proximal term** $(\gamma/2) \|\mathbf{x}_1 - \mathbf{x}_1^k\|_2^2$ is **neglected**.

Enhancements

Update the parameter κ

- ▶ **Constant step-size:** We can **fix** $\kappa_k = \kappa > 0$.
- ▶ **Increasing step-size:** κ_k can be **increased** as $\kappa_{k+1} := \eta\kappa_k$, for $k \geq 0$ and $\eta > 1$.
- ▶ **Adaptive step size:** κ_k can be updated **adaptively** based on **the primal and dual residuals** (see [2]).

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- ▶ **Adaptive step size:** κ_k can be updated **adaptively** based on **the primal and dual residuals** (see [2]).

Preconditioned ADMM

- ▶ **Drawback:** When \mathcal{X}_1 and \mathcal{X}_2 are **absent**, f_1 and f_2 possess a **tractable prox-operator**, if \mathbf{A}_1 and \mathbf{A}_2 are **not column orthogonal**, then we can not exploit the **proximal tractability** of f_1 and f_2 .
- ▶ **Overcome:** **Linearize** the quadratic terms and using the **gradient step** to **approximate** \mathbf{x}_1^{k+1} and \mathbf{x}_2^{k+1} :

$$\left\{ \begin{array}{ll} \mathbf{g}_1^k & := \mathbf{x}_1^k - \alpha_k^1 \mathbf{A}_1^T (\mathbf{A}_1 \mathbf{x}_1^k + \mathbf{A}_2 \mathbf{x}_2^k - \mathbf{b}) & \text{(gradient step for } \mathbf{x}_1) \\ \mathbf{x}_1^{k+1} & := \text{prox}_{\alpha_k^1 \kappa^{-1} f_1} (\mathbf{g}_1^k + \kappa^{-1} \mathbf{A}_1^T \lambda^k) & \text{(proximal step for } \mathbf{x}_1) \\ \mathbf{g}_2^k & := \mathbf{x}_2^k - \alpha_k^2 \mathbf{A}_2^T (\mathbf{A}_1 \mathbf{x}_1^{k+1} + \mathbf{A}_2 \mathbf{x}_2^k - \mathbf{b}) & \text{(gradient step for } \mathbf{x}_2) \\ \mathbf{x}_2^{k+1} & := \text{prox}_{\alpha_k^2 \kappa^{-1} f_2} (\mathbf{g}_2^k + \kappa^{-1} \mathbf{A}_2^T \lambda^k) & \text{(proximal step for } \mathbf{x}_2). \end{array} \right.$$

where α_k^1 and α_k^2 can be chosen **proportionally** to $\|\mathbf{A}_1\|^2$ and $\|\mathbf{A}_2\|^2$, respectively.

Convergence of ADMM

Theorem (Convergence of ADMM [2])

Assume that f_1 and f_2 are *proper, closed and convex* and \mathcal{L} has a *saddle point* $(\mathbf{x}^*, \lambda^*)$. For $\gamma = 0$, we have

- ▶ **Residual convergence:** $\{r_k\}$ converges to zero, where

$$r_k := \|\mathbf{A}_1 \mathbf{x}_1^k + \mathbf{A}_2 \mathbf{x}_2^k - \mathbf{b}\|_2.$$

- ▶ **Objective convergence:** $\{f(\mathbf{x}^k)\}$ converges to f^* .
- ▶ **Dual variable convergence:** $\{\lambda^k\}$ converges to λ^* .

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- ▶ **Dual variable convergence:** $\{\lambda^k\}$ converges to λ^* .

Theorem (Convergence rate of ADMM [4])

Let $\{\mathbf{w}^k\}$ be the sequence generated by ADMM, where $\mathbf{w}^k := (\mathbf{x}^k, \lambda^k)$ and $\mathbf{w}^* := (\mathbf{x}^*, \lambda^*)$. Let $\bar{\mathbf{w}}^k := (k+1)^{-1} \sum_{j=0}^k \mathbf{w}^j$. Then $\{\bar{\mathbf{w}}^k\}$ satisfies

$$f(\bar{\mathbf{x}}^k) - f(\mathbf{x}^*) + (\bar{\mathbf{w}}^k - \mathbf{w}^*)^T M(\mathbf{w}^*) \leq \frac{1}{2(k+1)} \|\mathbf{w}^0 - \mathbf{w}^*\|_{\mathbf{H}}^2, \quad \forall k \geq 0,$$

where $M(\mathbf{w}) := \begin{bmatrix} -\mathbf{A}^T \lambda \\ \mathbf{A}_1 \mathbf{x}_1 + \mathbf{A}_2 \mathbf{x}_2 - \mathbf{b} \end{bmatrix}$ and $\mathbf{H} := \text{diag}(\sqrt{\gamma} \mathbf{I}, \kappa \mathbf{A}_2^T \mathbf{A}_2, \kappa^{-1} \mathbf{I})$.

Consequently, $\{\mathbf{w}^k\}$ converges to \mathbf{w}^* at $\mathcal{O}(1/k)$ rate.

Example 1: Robust principle component analysis (RPCA)

Robust PCA

$$\begin{aligned} \min_{\mathbf{L}, \mathbf{S}} \quad & \|\text{vec}(\mathbf{S})\|_1 + \rho \|\mathbf{L}\|_*, \\ \text{s.t.} \quad & \mathbf{S} + \mathbf{L} = \mathbf{M}. \end{aligned} \tag{29}$$

Here $\rho > 0$ is a **weighted parameter** between the **sparse** and **low-rank** terms.

Example 1: Robust principle component analysis (RPCA)

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Here $\rho > 0$ is a **weighted parameter** between the **sparse** and **low-rank** terms.

Applying ADMM

The main steps of ADMM applying to (29) become:

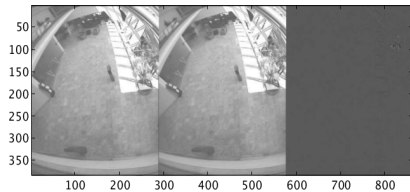
$$\begin{cases} \mathbf{S}^{k+1} & := \text{prox}_{\kappa^{-1} \|\text{vec}(\cdot)\|_1} (\mathbf{M} - \mathbf{L}^k + \kappa^{-1} \mathbf{W}^k), \\ \mathbf{L}^{k+1} & := \text{prox}_{\beta \kappa^{-1} \|\cdot\|_*} (\mathbf{M} - \mathbf{S}^{k+1} + \kappa^{-1} \mathbf{W}^k), \\ \mathbf{W}^{k+1} & := \mathbf{W}^k + \kappa (\mathbf{S}^k + \mathbf{L}^k - \mathbf{M}). \end{cases}$$

These **prox-operators** are computed as

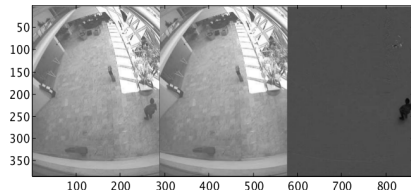
$$\begin{aligned} \text{prox}_{\tau \|\text{vec}(\cdot)\|_1}(\mathbf{S}) &= \text{sign}(\mathbf{S}_1) \otimes \max\{|\mathbf{S}_1| - \tau, 0\}, \\ \text{prox}_{\tau \|\cdot\|_*}(\mathbf{L}) &= \mathbf{U} \Sigma_\tau \mathbf{V}^T, \end{aligned}$$

where $\Sigma_\tau := \text{sign}(\Sigma) \otimes \max\{|\Sigma| - \tau, 0\}$ and $\mathbf{U} \Sigma \mathbf{V}^T = \mathbf{L}$ is the **SVD factorization** of \mathbf{L} .

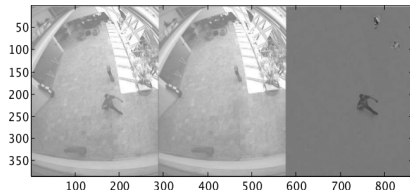
Video surveillance



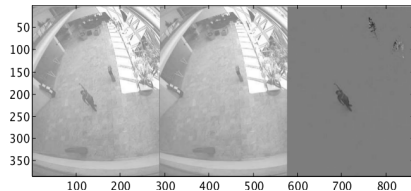
Frame 1



Frame 34



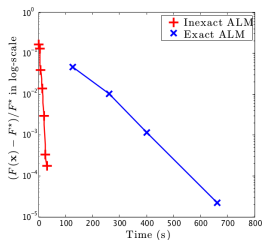
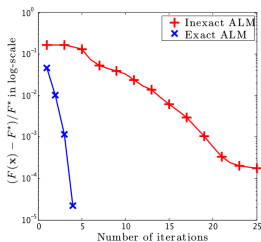
Frame 67



Frame 100

Unprocessed video from EC Funded CAVIAR project/IST 2001 37540, homepages.inf.ed.ac.uk/rbf/CAVIAR/.

Numerical test



	Exact ALM	Inexact ALM
Objective Value	553.5×10^3	553.6×10^3
Feas. Gap	0.33×10^{-5}	0.45×10^{-5}
$\ \mathbf{L}\ _*$	474.9×10^3	471.1×10^3
$\ \text{vec}(\mathbf{S})\ _1$	22.4616×10^6	23.556×10^6
Iterations	5	25
CPU time (s)	719.7	32.7
SVD Operations	644	25
Rank	1	1
Sparsity (%)	19.3	20.5

Algorithm

► Input

- M is 110592×100 : 100 frames of 288×384 pixels as columns

► Algorithm

- $\rho = 0.35 \times 10^{-2}$ - tunneable
- Stopping criteria: $\|M - L^k - S^k\| < 10^{-5} \|M\|$

Exact ADMM

Inexact ADMM

- (tunable)

$$\kappa^1 = 0.5 / \max\{\Sigma\}$$

$$\kappa^1 = 1.5 / \max\{\Sigma\}$$

- (tunable)

$$\kappa^{k+1} = \kappa^k * 6$$

$$\kappa^{k+1} = \kappa^k * 1.5$$

- prox op.

$$\text{Tolerance: } 10^{-6} \|M\|$$

Iterations: 1

► Output

- Numerical rounding \Rightarrow threshold
- $L_{\text{output}} = U_{\Sigma_{0.01 \max\{\Sigma\}}} V^T$
- $S_{\text{output}} = S_{0.01 \max\{|\Sigma|\}}$

Codes available at perception.csl.illinois.edu/matrix-rank/home.html

Example 2: Image deblurring

Image deblurring

The [image deblurring](#) presented previously can be written as:

$$\begin{array}{ll}
 \min_{\mathbf{u} \in \mathbb{R}^n \times p, \mathbf{v}} & \left\{ (1/2) \|\mathbf{v}\|_F^2 + \rho \|\mathbf{u}\|_{\text{TV}} \right\} \\
 \text{s.t.} & \mathcal{A}(\mathbf{u}) - \mathbf{v} = \mathbf{b}.
 \end{array} \tag{30}$$

Example 2: Image deblurring

Image deblurring

The **image deblurring** presented previously can be written as:

$$\begin{aligned} \min_{\mathbf{u} \in \mathbb{R}^n \times p, \mathbf{v}} \quad & \left\{ (1/2) \|\mathbf{v}\|_F^2 + \rho \|\mathbf{u}\|_{\text{TV}} \right\} \\ \text{s.t.} \quad & \mathcal{A}(\mathbf{u}) - \mathbf{v} = \mathbf{b}. \end{aligned} \quad (30)$$

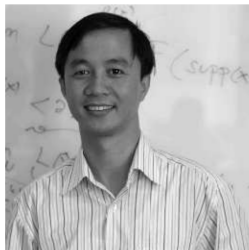
Applying ADMM

- ▶ We assume that $\mathcal{A}^* \mathcal{A} = \mathbf{I}$, where \mathcal{A}^* is the adjoint operator of \mathcal{A} .
- ▶ The \mathbf{v} -step can be computed **explicitly** and the \mathbf{u} -step can be computed relying on the **prox-operator** of the TV-norm.
- ▶ The main steps of ADMM becomes

$$\begin{cases} \mathbf{v}^{k+1} & := (\kappa + 1)^{-1} (\lambda^k + \kappa (\mathcal{A}(\mathbf{u}^k) - \mathbf{b})), \\ \mathbf{u}^{k+1} & := \text{prox}_{\rho \kappa^{-1} \|\cdot\|_{\text{TV}}} (\mathcal{A}^* (\mathbf{b} + \mathbf{v}^{k+1}) - \kappa^{-1} \lambda^k), \\ \lambda^{k+1} & := \lambda^k + \kappa (\mathcal{A}(\mathbf{u}^{k+1}) - \mathbf{v}^{k+1} - \mathbf{b}). \end{cases}$$

Wrong regularization parameter

$$\rho = \pi^e$$



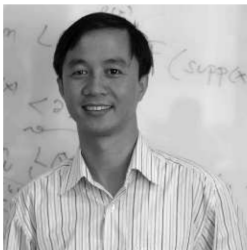
Original image



Blurred image
SNR = 40dB

Wrong regularization parameter

$$\rho = \pi^e$$



Original image

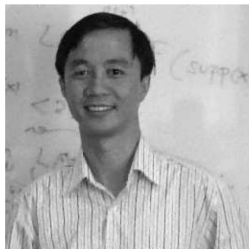


Blurred image
SNR = 40dB

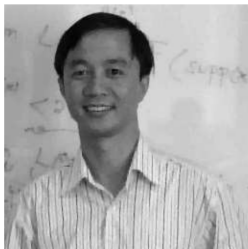


Recoverd image

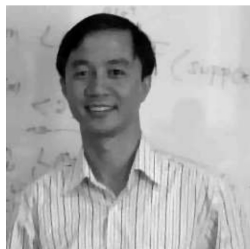
Different values of regularization parameter



$$\rho = 5 \times 10^{-3}$$



$$\rho = 1 \times 10^{-2}$$



$$\rho = 2.5 \times 10^{-2}$$

Numerical results

	$\rho = 5 \times 10^{-3}$	$\rho = 1 \times 10^{-2}$	$\rho = 2.5 \times 10^{-2}$
Objective Value	5317	7600	13344
MSE	24.1	22.8	27.2
ISNR (dB)	7.73	7.97	7.2
Feas. Gap ($\times 10^{-4}$)	3.01	3.38	5.45
Iterations	48	47	37
CPU time (s)	3.46	3.24	2.59
Linear Op. Calls*	99	97	77

► Algorithm

- $\kappa = \rho/10$
- Stopping criteria: $|F(\mathbf{u}^k, \mathbf{v}^k) - F(\mathbf{u}^{k-1}, \mathbf{v}^{k-1})| < 10^{-5} F(\mathbf{u}^k, \mathbf{v}^k)$
- Maximum 5 iterations for TV prox-operator (with warmstart)
- Input: 256px \times 256px image

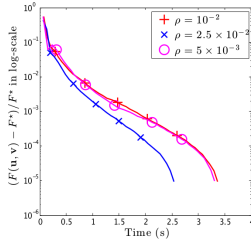
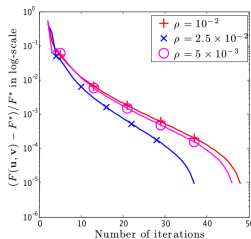
► MSE (Mean Squared Error) = $\frac{\|\mathbf{u} - \mathbf{u}^{\hat{h}}\|_2}{np}$

► ISNR (Improvement in Signal-to-Noise Ratio) = $\frac{\|\mathbf{b} - \mathbf{u}^{\hat{h}}\|_2}{np \text{MSE}} [\text{dB}]$

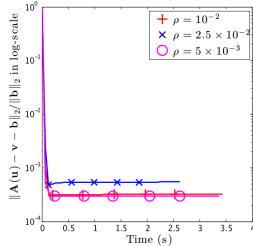
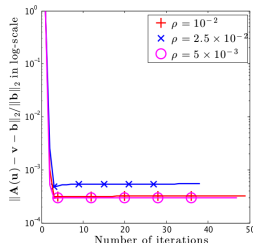
* number of applications of \mathbf{A} and \mathbf{A}^T operators

Convergence plots

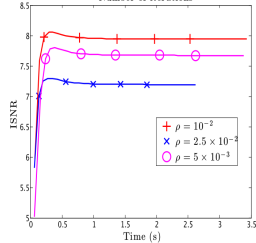
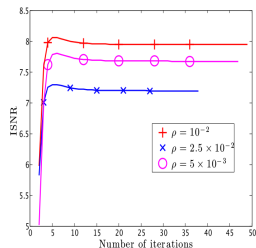
Objective



Feasibility Gap



ISNR

ISNR₀ = -20dB

Summary

We have studied **several methods** for solving the following **constrained convex problem**:

$$f^* := \min_{\mathbf{x}} \{f(\mathbf{x}) : \mathbf{Ax} = \mathbf{b}, \mathbf{x} \in \mathcal{X}\}. \quad (1)$$

Under **different assumptions**, we have presented the following methods:

- ▶ Null-space, projected gradient and Frank-Wolf's methods.
- ▶ Dual subgradient and augmented Lagrangian methods
- ▶ Alternating minimization algorithm (AMA) and alternating direction methods of multipliers (ADMM).

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However, such methods still have **limitations**, few of them are listed below.

Methods	Limitations
Null-space method	require null-space representation (e.g., QR with $\mathcal{O}(n^2 p)$ complexity), destroy the original structure of f
Projected gradient	require tractability of the projection on \mathcal{X} , smooth f
Dual subgradient method	advantage for decomposable structure , but slow convergence rate $\mathcal{O}(1/\sqrt{k})$, sensitive with the choices of step-size
Augmented Lagrangian	non-separability of the quadratic term, high-computational cost for subproblems , no supporting theory for penalty parameter selection
AMA	only application for partly strongly convex objective , not using the tractable proximity of f due to linear operator, no supporting theory for penalty parameter selection
ADMM	not using the tractable proximity of f due to linear operator, no supporting theory for penalty parameter selection

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In the **next lecture**, we will present **other methods** for solving (1) that either use **different set of assumptions** or **overcome some of these limitations**.

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